Structure in mappings on the random graph

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- Groups containing Aut(G)
- 3 Monoids containing Aut(G)
- 4 Model-theoretic corollaries
- 5 Ramsey theoretic tools

Denote by G = (V; E) the *random graph*, i.e., the unique graph which is

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Is this exceptional?

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General statement: Countable homogeneous structures are exactly the Fraïssé limits of Fraïssé classes of finite structures.

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Examples: Homogeneous K_n -free graph, dense linear order.

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How many classes of *n*-tuples are there?

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Do all homogeneous structures have oligomorphic automorphism groups?

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All homogeneous structures in a finite language have oligomorphic automorphism groups.

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What can we say about structures with comparable automorphism groups?

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Let Γ be homogeneous. Let Δ be any structure. Then Δ has a first-order definition in Γ iff $Aut(\Delta)$ contains $Aut(\Gamma)$.

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Other examples

There are 5 groups containing Aut((\mathbb{Q} , <)) (Cameron '76). There are 2 groups containing the homogeneous K_n -free graph (Thomas '91 / P '09).

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Is the number of groups containing the automorphisms of a homogeneous structure always finite?

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Conjecture (Thomas '91)

The number of groups containing the automorphisms of a homogeneous structure is always finite.

M. Pinsker (Caen)

Mappings on the random graph

Why did you reprove a 20-years old theorem?

Why?

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Definability

Groups: First-order interdefinability. Monoids: Existential positive interdefinability. Clones: Primitive positive interdefinability.

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Corollary

Let Γ be a reduct of the random graph. Then Γ is a reduct of (V; =), or its endomorphisms are generated by its automorphisms.

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What does this have to do with mappings on G?

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Theorem (Bodirsky & P '09)

All reducts of the random graph are model-complete.

Ramsey's theorem

M. Pinsker (Caen)

Ramsey's theorem

Let
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Given h, p, can we choose n large enough such that $n \rightarrow (h)^p$ holds?

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Theorem (Ramsey's theorem)

For all p, h there exists n such that $n \to (h)^p$.

M. Pinsker (Caen)

Let *N*, *H*, *P* be graphs.

$$N \to (H)^P$$

means:

For all partitions of the copies of P in N into good and bad there exists a copy of H in Nsuch that the copies of P in H are all good or all bad. Let *N*, *H*, *P* be graphs.

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means:

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Definition

A class $\ensuremath{\mathbb{C}}$ of structures of the same signature is called a Ramsey class iff

for all $H, P \in \mathbb{C}$ there is N in \mathbb{C} such that $N \to (H)^P$.

Theorem (Nešetřil-Rödl)

The set of finite ordered graphs is a Ramsey class.

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Conjecture

Any homogeneous structure **whose set of finite induced substructures (+ order) is a Ramsey class** has only finitely many reducts.

Enjoy your coffee break!

Where can we find your paper?

Manuel Bodirsky and Michael Pinsker,

All reducts of the random graph are model-complete, available from arXiv.