

# *Abstract Elementary Classes and their axiomatizations: a review*

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# Abstract Elementary Classes and their axiomatizations: a review

- Shelah's Presentation Theorem
- Relational Presentation Theorem
- Leung's Axiomatization
- The semantic-syntactic correspondence
- Shelah-Villaveces Axiomatization
- What we have

# AEC definition

Let  $\tau$  a language and  $\mathcal{K}$  a class of  $\tau$ -structures. We say that  $(\mathcal{K}, \prec_{\mathcal{K}})$  is an Abstract Elementary Class (AEC) if and only if:

- ①  $\prec_{\mathcal{K}}$  is a partial order over  $\mathcal{K}$  and refines the  $\subseteq$  relation.
- ②  $(\mathcal{K}, \prec_{\mathcal{K}})$  is closed under isomorphisms.
- ③ There is a Cardinal  $\kappa$ , called the Löwenheim-Skolem number of  $\mathcal{K}$ , such that for all  $M \in \mathcal{K}$  and all  $A \subseteq |M|$  there is  $N \in \mathcal{K}$  such that  $A \subseteq |N|$  and  $\|N\| \geq \kappa$ .
- ④ For all  $M_1, M_2, N \in \mathcal{K}$  such that  $M_1, M_2 \prec_{\mathcal{K}} N$  and  $M_1 \subseteq M_2$ , then  $M_1 \prec_{\mathcal{K}} M_2$ .
- ⑤ For every increasing and continuous  $\prec_{\mathcal{K}}$ -chain  $\langle M_i \rangle_{i < \alpha}$ , we have:
  - $M_{\alpha} := \bigcup_{i < \alpha} M_i \in \mathcal{K}$ .
  - For all  $i < \alpha$ ,  $M_i \prec_{\mathcal{K}} \bigcup_{i < \alpha} M_i$ .
  - If  $N \in \mathcal{K}$  is such that  $M_i \prec_{\mathcal{K}} N$ , then  $M_{\alpha} \prec_{\mathcal{K}} N$ .

## Example

- If  $T$  is a first order theory, then  $(Mod(T), \prec_{\mathcal{K}})$  is an AEC with  $LS(\mathcal{K}) = \aleph_0$ .
- If  $\psi \in \mathbb{L}_{\omega_1, \omega}$  and  $\Delta \subseteq \mathbb{L}_{\omega_1, \omega}$  a countable fragment that contains  $\psi$ , then  $(Mod(\psi), \prec_{\Delta})$  is an AEC with  $LS(\mathcal{K}) = \aleph_0$ .
- (Mazari-Armida 23) Classes of abelian groups and modules using the pure subgroup or module relation.

## Definition

Let  $f : \mathcal{M} \longrightarrow \mathcal{N}$  be an embedding,  $f$  is a  $\mathcal{K}$ -embedding if  $f[\mathcal{M}] \prec_{\mathcal{K}} \mathcal{N}$

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# Shelah's Presentation Theorem and the EM-functor

## Theorem (Sh:87)

Let  $K = (\mathcal{K}, \prec_K)$  be an AEC with Löwenheim-Skolem number  $\kappa$  in a vocabulary  $\tau$  such that  $|\tau| \leq \kappa$ . Then, there exists a vocabulary  $\tau' \supset \tau$  with  $|\tau'| = \kappa$ , a first order  $\tau'$ -theory  $T'$  and a set  $\Gamma'$  of quantifier free  $T'$ -types such that

$$\mathcal{K} = \{M' \restriction_{\tau} : M' \models T' \text{ and omits all the types in } \Gamma'\}.$$

Furthermore,

- ① If  $M', N' \models T'$  are such that that omit all the types in  $\Gamma'$  and  $M' \subseteq N'$ , then  $M' \restriction_{\tau} \prec_K N' \restriction_{\tau}$  and,
- ② If  $M, N \in K$  are such that  $M \prec_K N$ , then there are expansions  $M'$  of  $M$  and  $N'$  of  $N$  to  $\tau'$  such that  $M', N' \models T'$ , omit all the types in  $\Gamma'$  and  $M' \subseteq N'$ .

$$(\mathcal{K}, \prec_K) = (PC_{\tau}(T', \Gamma'), \subseteq)$$

# Extracting indiscernibles

**Definition 11.5.** Let  $\Phi$  be an EM blueprint. Let  $I, J$  be linear orders, let  $\delta$  be a limit ordinal and let  $\langle \bar{a}_j : j \in J \rangle$  be a sequence. We say that  $\langle \bar{a}_j : j \in J \rangle$  is  $(\Phi, I)$ -strictly indiscernible if:

- (1)  $J$  is infinite.
- (2) For some  $\alpha$ , for all  $j \in J$ ,  $\bar{a}_j \in {}^\alpha \text{EM}_\tau(I, \Phi)$ .
- (3) There exists a sequence  $\langle \bar{a}'_j : j \in J \rangle$  and a sequence of terms  $\bar{\rho}$  such that  $\bar{a}_j = \bar{\rho}(\bar{a}'_j)$  for all  $j \in J$  and  $\langle \bar{a}'_j : j \in J \rangle$  is quantifier-free indiscernible in the vocabulary of linear orders inside  $I$ .

We call  $\langle \bar{a}_j : j \in J \rangle$   $(\Phi, I)$ -strictly indiscernible over  $A$  if  $\langle \bar{a}_j \bar{a} : j \in J \rangle$  is  $(\Phi, I)$ -strictly indiscernible for some (any) enumeration  $\bar{a}$  of  $A$ .

**Theorem 11.7** (Strict indiscernible extraction). Let  $\mathbf{K}$  be an AEC with arbitrarily large models and let  $\text{LS}(\mathbf{K}) < \theta \leq \lambda$  be cardinals with  $\theta$  regular. Let  $\kappa < \theta$  be a (possibly finite) cardinal. Let  $\Phi \in \Upsilon_{\text{LS}(\mathbf{K})}[\mathbf{K}]$  be an EM blueprint for  $\mathbf{K}$ .

Let  $N := \text{EM}_\tau(\mathbf{K})(\lambda, \Phi)$ . Let  $M \in \mathbf{K}_{\leq \text{LS}(\mathbf{K})}$  be such that  $M \leq_{\mathbf{K}} N$ . Let  $\langle \bar{a}_i : i < \theta \rangle$  be a sequence of distinct elements such that for all  $i < \theta$ ,  $\bar{a}_i \in {}^\kappa |N|$ .

If  $\theta_0^\kappa < \theta$  for all  $\theta_0 < \theta$ , then there exists  $w \subseteq \theta$  with  $|w| = \theta$  such that  $\langle \bar{a}_i : i \in w \rangle$  is  $(\Phi, \lambda)$ -strictly indiscernible over  $M$ .

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# Relational Presentation Theorem

## Theorem (BaBo17)

Let  $(\mathcal{K}, \prec_{\mathcal{K}})$  be an AEC with Löwenheim-Skolem number  $\kappa$  in a vocabulary  $\tau$ . Then there exists an expansion of  $\tau$  by predicates of arity  $\kappa$  and a  $T'$   $\tau'$ -theory in  $\mathbb{L}_{(2^{\kappa})^{+}, \kappa^{+}}$  such that

$$\mathcal{K} = \{M' \upharpoonright_{\tau} : M' \models T'\}.$$

Furthermore

- ① If  $M', N' \models T'$  are such that  $M' \subseteq N'$ , then  $M' \upharpoonright_{\tau} \prec_{\mathcal{K}} N' \upharpoonright_{\tau}$ .
- ② If  $M, N \in \mathcal{K}$  are such that  $M \prec_{\mathcal{K}} N$ , then there are expansions  $M'$  of  $M$  and  $N'$  of  $N$  to  $\tau'$  such that  $M', N' \models T'$  and  $M' \subseteq N'$ .

$$(\mathcal{K}, \prec_{\mathcal{K}}) = (PC_{\tau}(T'), \subseteq)$$

## Theorem (BaBo17)

Let  $\kappa$  be a strongly compact cardinal and let  $(\mathcal{K}, \prec_{\mathcal{K}})$  be an AEC with  $LS(\mathcal{K}) < \kappa$ . If  $\mathcal{K}_{\leq \mu, \kappa}$  has AP, JEP..., then  $\mathcal{K}_{> \mu}$  has AP, JEP....

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# Leung's Axiomatization

## Theorem (Le23)

*Let  $\lambda = \kappa + I_2(\kappa, \mathcal{K})$  where  $I_2(\kappa, \mathcal{K})$  is the number of non-isomorphic pairs  $(M, N)$  such that  $M \prec_{\mathcal{K}} N$  and bought have cardinality  $LS(\mathcal{K})$ . There is  $\sigma_{\mathcal{K}} \in \mathbb{L}_{\lambda^+, \kappa^+}(\omega \cdot \omega)(\tau)$  such that  $(\mathcal{K}, \prec_{\mathcal{K}}) = (\{M \in \tau\text{-structures} \mid M \models \sigma_{\mathcal{K}}\}, \prec_{\Delta})$ .*

## Remark

*This is used to simplify some resoults and extend it to other contexts.*

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# The correspondence

## Definition (Galois Morlization, Vas16)

Let  $\kappa$  be an infinite cardinal and let  $(\mathcal{K}, \prec_{\mathcal{K}})$  be an AEC. The  $(< \kappa)$ -Galois Morlization of  $(\mathcal{K}, \prec_{\mathcal{K}})$  is  $(\hat{\mathcal{K}}, \prec_{\hat{\mathcal{K}}})$ , an AEC in a  $< \kappa$ -ary language  $\hat{\tau}$  extending  $\tau$  such that:

- ①  $\mathcal{K} = \hat{\mathcal{K}}$  and  $\prec_{\mathcal{K}} = \prec_{\hat{\mathcal{K}}}$ .
- ② For all  $p \in ga-S^{<\kappa}(\emptyset; \mathbb{C})$ , there exists  $R_p \in \hat{\tau}$  such that  $\mathbb{C} \models R_p[\bar{b}]$  iff  $p = ga-tp(\bar{a}/\emptyset; \mathbb{C})$ .
- ③  $tp_{\Delta}(\bar{b}/A; \mathbb{C}) := \{\psi(\bar{x}; \bar{a}) \in qf-\mathbb{L}_{\kappa, \kappa}(\hat{\tau}) \text{ formulas} \mid \bar{a} \in A \text{ y } \mathbb{C} \models \psi[\bar{b}; \bar{a}]\}$ .
- ④  $qf-\mathbb{L}_{\kappa, \kappa}-S^{<\kappa}(A; \mathbb{C}) := \{tp_{\Delta}(\bar{b}/A; \mathbb{C}) \mid \bar{b} \in |\mathbb{C}|\}$ .

## Fact (The semantic-syntactic correspondence, Vas16)

$\mathcal{K}$  is  $(< \kappa)$ -tame iff  $ga-tp(\bar{b}/A; \mathbb{C}) \mapsto tp_{\Delta}(\bar{b}/A; \mathbb{C})$  from  $ga-S^{<\kappa}(A; \mathbb{C})$  to  $qf-\mathbb{L}_{\kappa, \kappa}-S^{<\kappa}(A; \mathbb{C})$  is a bijection.

# Order Property

## Definition

We say that  $\mathcal{K}$  has the  $(\kappa_1, \kappa_2, \theta)$ -order property of length  $\mu$  if there are  $A \subseteq |\mathbb{C}|$  with  $|A| \leq \theta$ ,  $\langle \bar{a}_i | i < \mu \rangle$  where  $\bar{a}_i \in^{\kappa_1} |\mathbb{C}|$  and  $\langle \bar{b}_i | i < \mu \rangle$  where  $\bar{b}_i \in^{\kappa_2} |\mathbb{C}|$  such that if  $i_0 < j_0 < \mu$ ,  $i_1 < j_1 < \mu$ , then  $ga\text{-}tp(\bar{a}_{i_0} \bar{b}_{j_0} / A; \mathbb{C}) \neq ga\text{-}tp(\bar{a}_{j_1} \bar{b}_{i_1} / A; \mathbb{C})$ .

## Fact

*If  $\mathcal{K}$  has AP and is  $\kappa$ -tame, then  $\mathcal{K}$  is  $\lambda$ -stable iff does not have the  $(\kappa_1, \kappa_2, \lambda)$ -order property.*

# Independence property

## Definition

Let  $\lambda$  be a cardinal.

$ded(\lambda) = \sup\{\kappa : \text{there is a linear order of size } \kappa \text{ which has a dense subset of size } \lambda\}.$

## Fact

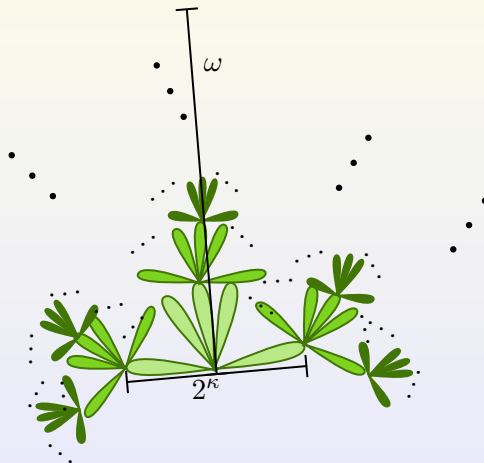
*Let  $\mu := \beth_{(2^{LS(\mathcal{K})})_+}$ . If  $\mathcal{K}$  is  $< \aleph_0$ -tame,  $C \subseteq |\mathbb{C}|$  with  $|C| = \lambda > \beth_3(LS(\mathcal{K}))$  and  $|ga-S^1(C; \mathbb{C})| > Ded(\lambda)$ , y  $Ded(\lambda)^{2^\kappa} = Ded(\lambda)$ , then there are  $\psi(\bar{x}, \bar{y}) \in qf - \mathbb{L}_{\kappa, \kappa}(\hat{\tau})$ ,  $\langle \bar{a}_i \in |\mathbb{C}| : i < \mu \rangle$  y  $\langle \bar{b}_w \in |\mathbb{C}| : w \subseteq \mu \rangle$  such that  $\mathbb{C} \models \psi[\bar{a}_i; \bar{b}_w]$  iff  $i \in w$ .*

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# Canonical tree



# The formulas

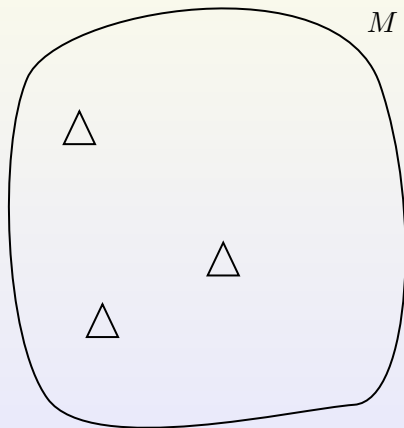
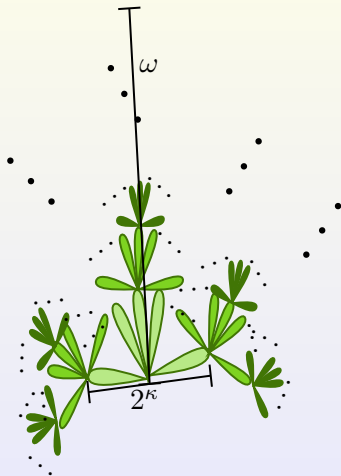
We define by induction on  $\gamma < \lambda^+$  formulas  $\varphi_{N,\gamma,n}(\bar{x}_n)$  in  $\mathbb{L}_{\lambda^+,\kappa^+}$  for all  $n < \omega$  and  $N \in \mathcal{S}_n$ .

- ① If  $\gamma = 0$ :
  - ① If  $n = 0$ , then  $\varphi_{\emptyset,0,0}(\bar{x}_n) := \top$ .
  - ② For  $n > 0$ , let  $\varphi_{N,0,n}(\bar{x}_n) := \bigwedge \text{Diag}_n^\kappa(N)$  where  $\text{Diag}_n^\kappa(N) := \{\varphi(x_{\alpha_0}, \dots, x_{\alpha_{k-1}}) : \alpha_0, \dots, \alpha_{k-1} < \kappa \cdot n, \varphi(x_{\alpha_0}, \dots, x_{\alpha_{k-1}}) \text{ is an atomic or negation of an atomic formula and } N \models \varphi(a_{\alpha_0}^*, \dots, a_{\alpha_{k-1}}^*)\}$ .
- ② If  $\gamma = \beta + 1$ , then

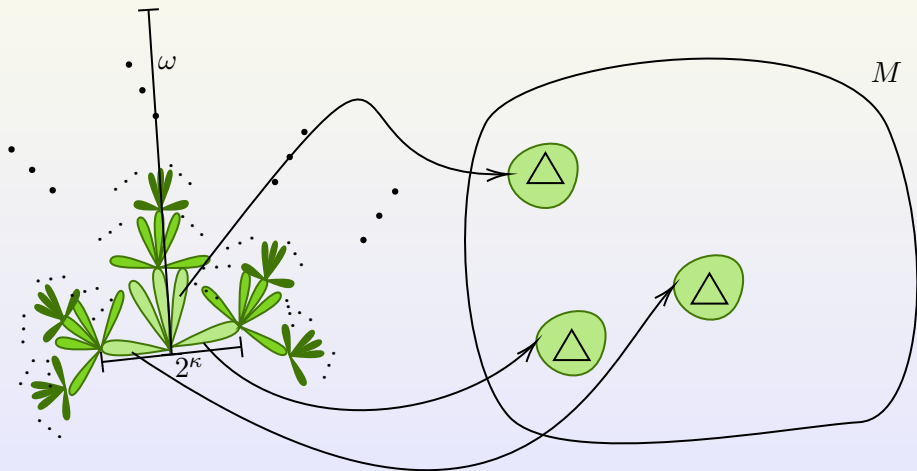
$$\varphi_{N,\gamma,n}(\bar{x}_n) := \bigvee_{N \prec_{\mathcal{K}} N', N' \in \mathcal{S}_{n+1}} \exists \bar{x}_{=n} \left[ \varphi_{N',\beta,n+1}(\bar{x}_{n+1}) \wedge \bigwedge_{\alpha < \kappa} \bigvee_{\delta < \kappa \cdot (n+1)} z_\alpha = x_\delta \right].$$

- ③ If  $\gamma$  is a limit ordinal, then  $\varphi_{N,\gamma,n}(\bar{x}_n) := \bigwedge_{\beta < \gamma} \varphi_{N,\beta,n}(\bar{x}_n)$

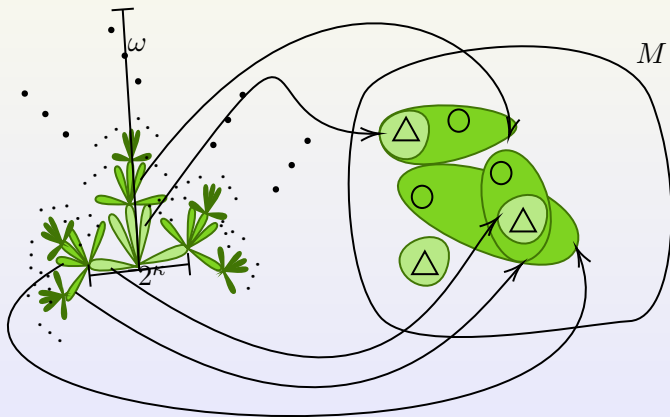
$$M \models \varphi_{\emptyset,1,0}$$



$$M \models \varphi_{\emptyset,1,0}$$



$$M \models \varphi_{\emptyset,2,0}$$



# The sentence

## Fact

Let  $\lambda = \beth_2(\kappa)^{++}$ . We have  $\varphi_{\emptyset, \lambda+1, 0} \in \mathbb{L}_{\lambda^+, \kappa^+}$ .

## Theorem (Shelah-Villaveces 2022)

$M \in \mathcal{K}$  if and only if  $M \models \varphi_{\emptyset, \lambda+1, 0}$

## Proof sketch.

Notice that  $M \models \varphi_{\emptyset, \lambda+1, 0}$  if and only if for all  $\gamma < \lambda$  and all  $A \in [|M|]^\kappa$ ,  $M \models \varphi_{N, \gamma, 1}[A]$  for some  $N \in \mathcal{S}_1$ .

Left to right: induction on  $\gamma < \lambda$ . Use coherence and Löwenheim-Skolem.

Right to left: show that

$\mathbb{S} := \{M^* \subseteq M \mid \text{there are } N \in \mathcal{S}_1 \text{ with enumeration } \langle a_\alpha^* \mid \alpha < \kappa \rangle \text{ and}$

$f : N \cong M^* \text{ such that } M \models \varphi_{N, \lambda, 1}[\langle f(a_\alpha^*) \mid \alpha < \kappa \rangle]\}$

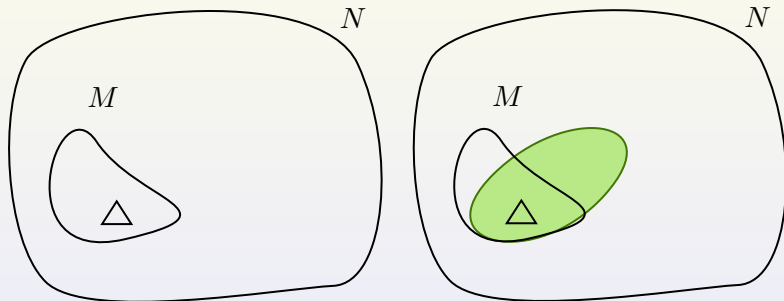
is a directed system. Use a complicate combinatorial principle. □

# Syntactic substructure criteria

Theorem (Shelah-Villaveces 2022, N.S. 2023)

- ①  $M_1 \prec_{\mathcal{K}} M_2$ ,
- ② if  $\bar{a} \in |M_1|^{\leq \kappa}$  then there are  $M_{\bar{a}} \prec_{\mathcal{K}} M_1$ ,  $N_{\bar{a}} \in \mathcal{S}_1$  with enumeration  $\langle a_{\alpha}^* | \alpha < \kappa \rangle$  and  $f_{\bar{a}} : N_{\bar{a}} \cong M_{\bar{a}}$  such that
  - ①  $\bar{a} \in |M_{\bar{a}}|^{\leq \kappa}$  and
  - ②  $M_2 \models \varphi_{N_{\bar{a}}, \lambda, 1}[\langle f_{\bar{a}}(a_{\alpha}^*) | \alpha < \kappa \rangle]$ .

# Syntactic substructure criteria



$$(\mathcal{K}, \prec_{\mathcal{K}}) = (\text{Mod}(\psi_{\mathcal{K}}), \prec_{\Delta})$$



# A game to know if $M \in \mathcal{K}$ : $G_{AEC}(M)$

Let  $M$  be a  $\tau(\mathcal{K})$ -structure. Remember that  $\lambda = \beth_2(\kappa)^{++}$  and  $\kappa = LS(\mathcal{K})$ . The states of the game are pairs  $(\alpha, \pi)$  where  $\alpha < \lambda$  and  $\pi : N \rightarrow M$  is a  $\mathcal{K}$ -embedding for  $N \in \mathcal{S}_n$  and  $n < \omega$ .

**Starting stage:** is  $(\lambda, \emptyset)$ .

**Further stages:** At stage  $(\alpha, \pi)$ :

- ① Player I: picks an ordinal  $\alpha < \lambda$  and a tuple  $\bar{a} \in |M|^\kappa$ .
- ② Player II: picks  $N' \in \mathcal{S}_{n+1}$  and a  $\mathcal{K}$ -embedding  $\pi' : N' \rightarrow M$  such that  $N \prec_{\mathcal{K}} N'$  and

$$\pi' = \pi \cup \{(n_i, m_i) : i \in [\kappa \cdot n, \kappa \cdot (n+1))\}$$

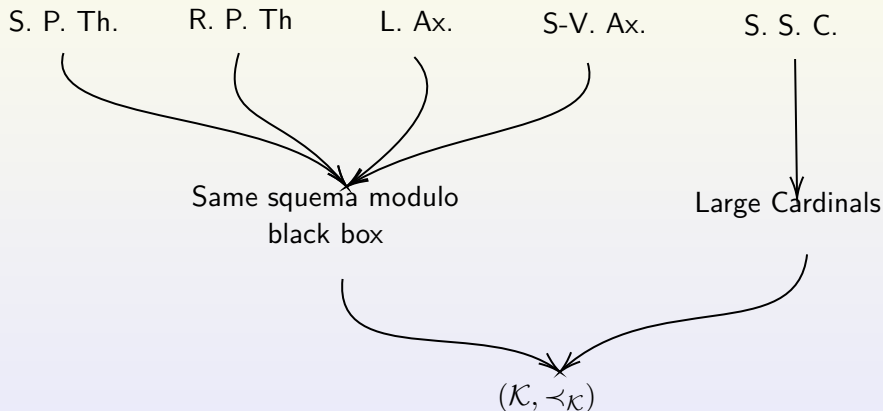
## Fact

$M \in \mathcal{K}$  iff player II has a winning strategy in the game  $G_{AEC}(M)$ .

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Thank you! :)