

# Clones (1&2)

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# Base set $X$

Let  $X$  be a (nonempty) set.

- ▶ Often finite:
  - ▶  $X = \{0, 1\}$ .
  - ▶  $X = \{0, *, 1\}$ .
  - ▶  $X = \{ \{\}, \{a\}, \{b\}, \{a, b\} \}$ .
  - ▶  $X = \{1, \dots, n\}$ .
  - ▶ Etc.
- ▶ Sometimes countably infinite:
  - ▶  $X = \mathbb{N} = \{0, 1, 2, \dots\}$ .
- ▶ Sometimes uncountably infinite:
  - ▶  $X = \mathbb{R}$ , etc.

# Operations on $X$

$X$  = our base set.

- ▶ A unary operation is a (total) function  $f : X \rightarrow X$ .
- ▶ A binary operation is a function  $f : X^2 \rightarrow X$ .
- ▶ ternary, quaternary, ...
- ▶ A  $k$ -ary operation is a function  $f : X^k \rightarrow X$  (for  $k \geq 1$ ).
- ▶ We write  $\mathcal{O}^{(k)}$  or  $\mathcal{O}_X^{(k)}$  for the set of all  $k$ -ary operations on  $X$ . (Sometimes also written  $X^{X^k}$ .)
- ▶ We let  $\mathcal{O}_X := \bigcup_{k=1}^{\infty} \mathcal{O}_X^{(k)}$ .

(For simplicity we will assume that the sets  $X^k$  are pairwise disjoint. We will ignore the 0-ary functions and replace them by constant 1-ary functions.)

# Transformation monoids

## Definition ((abstract) monoid)

A *monoid* or **abstract monoid** is a structure  $(M, *, 1)$ , where

- $*$  is a binary operation on  $M$ , associative
- ... together with a neutral element  $1$  ( $1 * a = a * 1 = a$ ).

## Definition (transformation/concrete monoid, unary clone)

A **transformation monoid** is a subset  $T \subseteq \mathcal{O}_X^{(1)}$  (for some  $X$ ) which is closed under composition and contains the identity function  $id : X \rightarrow X$ . ( $(T, \circ, id)$  will be an abstract monoid.)

Conversely, a variant of Cayley's theorem shows that every abstract monoid is isomorphic to a transformation monoid.

## Binary clones

A *transformation monoid* or *unary clone* on  $X$  is a subset  $T \subseteq \mathcal{O}_X^{(1)}$  which is closed under composition and contains the identity function  $id : X \rightarrow X$ .

### Definition

A **binary clone** on  $X$  is a set  $T \subseteq \mathcal{O}_X^{(1)}$  which is closed under "composition" and contains the two projections  $\pi_1, \pi_2 : X^2 \rightarrow X$ .

### Definition (Composition)

Let  $f, g_1, g_2 \in \mathcal{O}_X^{(2)}$ . The composition  $f(g_1, g_2)$  is the function from  $X^2$  to  $X$  defined by

$$f(g_1, g_2)(x, y) := f(g_1(x, y), g_2(x, y))$$

## $k$ -ary clones

### Definition ( $k$ -ary clone)

A  $k$ -ary clone on  $X$  is a set  $T \subseteq \mathcal{O}_X^{(k)}$  which is closed under "composition" and contains the  $k$  projections  $\pi_1, \dots, \pi_k : X^k \rightarrow X$ .

### Definition (Composition)

Let  $f, g_1, \dots, g_k \in \mathcal{O}_X^{(k)}$ . The composition  $f(g_1, \dots, g_k)$  is the function from  $X^k$  to  $X$  defined by

$$\forall \vec{x} \in X^k : f(g_1, \dots, g_k)(\vec{x}) := f(g_1(\vec{x}), \dots, g_k(\vec{x}))$$

("Plugging  $g_1, \dots, g_k$  into  $f$ ")

# Clones

## Definition (Clone)

A **clone** on  $X$  is a set  $T \subseteq \mathcal{O}_X = \bigcup_{k=1}^{\infty} \mathcal{O}_X^{(k)}$  which is closed under "composition" and contains all projections  $\pi_k^n : X^n \rightarrow X$ ,  $n = 1, 2, \dots$ ,  $1 \leq k \leq n$ .

## Definition (Composition)

Let  $f \in \mathcal{O}_X^{(k)}$ ,  $g_1, \dots, g_k \in \mathcal{O}_X^{(m)}$ . The composition  $f(g_1, \dots, g_k)$  is the function from  $X^m$  to  $X$  defined by

$$\forall \vec{x} \in X^m : f(g_1, \dots, g_k)(\vec{x}) := f(g_1(\vec{x}), \dots, g_k(\vec{x}))$$

("Plugging  $g_1, \dots, g_k$  into  $f$ ")

If  $C$  is a clone, then  $C^{(k)} := C \cap \mathcal{O}_X^{(k)}$  is a  $k$ -ary clone, the  **$k$ -ary fragment** of  $C$ .

## Examples of clones

- ▶ The smallest clone  $J_X$  contains only the projections.
- ▶ The largest clone  $\mathcal{O}_X$  contains all operations.
- ▶ Every subset  $S \subseteq \mathcal{O}_X$  will *generate* a clone  $\langle S \rangle$ , the smallest clone containing  $S$ . The clone  $\langle S \rangle$  can be obtained **from below** by closing  $S$  under composition, or **from above** as  $\langle S \rangle = \bigcap \{ M \mid S \subseteq M \subseteq \mathcal{O}_X, M \text{ is a clone} \}$ .
- ▶ If  $V$  is a vector space over the field  $K$ , then the set of all linear functions  $f_{\vec{a}} : V^k \rightarrow V$

$$f_{\vec{a}}(v_1, \dots, v_k) := a_1 v_1 + \dots + a_k v_k$$

(with  $\vec{a} = (a_1, \dots, a_k) \in K^k$ ) is a clone.



## Examples of clones, continued

For every algebra  $\mathcal{X} = (X, f, g, \dots)$  (=universe  $X$  with operations  $f, g, \dots$  — for example  $\mathcal{X}$  might be a group, a ring, etc) we consider

- ▶ the clone of *term operations* on  $X$ , the smallest clone containing all the basic operations  $f, g, \dots$  of  $\mathcal{X}$ ;
- ▶ the clone of *polynomial operations* on  $X$ , the smallest clone containing all terms as well as all constant unary functions on  $X$ .

Many properties of the algebra  $\mathcal{X}$  depend only on the clone of term functions, and not on the specific set of basic operations which generates this clone. (E.g. subalgebras, congruence relations, automorphisms, etc)

For example, a **Boolean algebra** will have the same clone as the corresponding **Boolean ring**.

# The family of all clones

For any nonempty set  $X$  let  $Cl(X)$  be the set of all clones on  $X$ .

- ▶ The intersection of any subfamily of  $Cl(X)$  is again in  $Cl(X)$ .
- ▶  $(Cl(X), \subseteq)$  is a complete lattice.  
Meet = intersection, join = generated by union.
- ▶  $J_X$  is the smallest clone,  $\mathcal{O}_X$  the largest.
- ▶ If  $X = \{0\}$ , then there is a unique clone:  $J_X = \mathcal{O}_X$ .
- ▶ If  $X = \{0, 1\}$ , then  $Cl(X)$  is countably infinite.
- ▶ If  $X$  is finite and has at least three elements, then  $Cl(X)$  is uncountable. (In fact:  $|Cl(X)| = |\mathbb{R}|$ .)
- ▶ If  $X$  is infinite, then ... (later)

# Uncountably many clones

If  $X = \{0, 1, 2\}$ , then  $Cl(X)$  is uncountable.

## Proof sketch.

- ▶ We call a  $k$ -tuple  $(a_1, \dots, a_k) \in \{0, 1, 2\}^k$  proper, if exactly one of the  $a_i$  is equal to 1, and all the others are 2.
- ▶ For every  $k \geq 3$  let  $f_k : X^k \rightarrow X$  be the function that assigns 1 to every proper  $k$ -tuple, and 0 to everything else.
- ▶ For every  $A \subseteq \{3, 4, \dots\}$  let  $C_A := \langle \{f_i \mid i \in A\} \rangle$ .
- ▶ Check that for  $k \notin A$  we have  $f_k \notin C_A$ .  
(Every composition of functions  $f_i, i \neq k$  will assign 0 to some proper  $k$ -tuple.)
- ▶ Hence the map  $A \mapsto C_A$  is 1-1.

# Completeness

Fix a base set  $X$ .

## Definition

A set  $S \subseteq \mathcal{O}_X$  is **complete** if  $\langle S \rangle = \mathcal{O}_X$ , i.e., if every operation on  $X$  is term function of the algebra with operations  $S$ .

## Example

Let  $X = \{0, 1\}$ ,  $\mathcal{X} = (X, \vee, \wedge, \neg, 0, 1)$ .

- ▶ The set  $\{\vee, \wedge, \neg\}$  is complete.
- ▶ The set  $\{\wedge, \neg\}$  is complete.
- ▶ The set  $\{\mid\}$  is complete, where  $x \mid y := \neg(x \wedge y)$ .  
(Sheffer stroke)

# Completeness, more examples

## Theorem

For every  $X$ :  $\langle \mathcal{O}_X^{(2)} \rangle = \mathcal{O}_X$ .

## Proof.

- ▶ finite: Lagrange interpolation
- ▶ infinite: use  $X \times X \approx X$ .

Caution: Most clones  $C$  are NOT generated by their binary fragment  $C \cap \mathcal{O}^{(2)}$ . (Not even finitely generated.)

## Theorem

If  $X = \{1, \dots, k\}$ , then there is a single function  $f \in \mathcal{O}_X^{(2)}$  with  $\langle f \rangle = \mathcal{O}_X^{(2)}$ : Let  $f(x, x) = x + 1$  (modulo  $k$ ),  $f(x, y) = 0$  otherwise.

## (Completeness on infinite sets)

If  $X$  is infinite, then  $\mathcal{O}_X$  is uncountable. Hence a finite/countable set of operations cannot generate all of  $\mathcal{O}_X$ .

However:

### Theorem

*Let  $X \neq \emptyset$ . For any finite or **countable** set  $T \subseteq \mathcal{O}_X$  there is a single function  $f_T$  (not necessarily in  $T$ ) such that  $T \subseteq \langle f \rangle$ .*

### Theorem

- If  $X$  is countable, then there is a **countable dense subset** of  $\mathcal{O}_X$  (in the natural topology), hence there is a single function  $f$  such that the **topological closure** of  $\langle f \rangle$  is all of  $\mathcal{O}_X$ .*
- If  $X$  is uncountable, then  $\mathcal{O}_X$  will **not** be **separable** any more.*

## Completeness, continued

Let  $X = \{0, 1\}$  be the 2-element Boolean algebra, with Boolean operations  $\wedge, \vee, \neg, \rightarrow, |, \dots$ .

### Example

The set  $\{\vee, \wedge, \rightarrow\}$  is not complete.

### Proof.

Each of the three operations **preserves** the set  $\{1\}$ , i.e., this set is a subalgebra of the algebra  $(\{0, 1\}, \wedge, \vee, \rightarrow)$ .

Hence every function in  $\langle\{\wedge, \vee, \rightarrow\}\rangle$  will also preserve this set, but  $\neg$  does not. So  $\neg \notin \langle\{\wedge, \vee, \rightarrow\}\rangle$ .

# Polymorphisms, example

## Example

The set  $\{\vee, \wedge, 0, 1\}$  is not complete.

## Proof.

All four functions are **monotone** in both arguments.

## Definition

Let  $\rho \subseteq X \times X$  be a relation (Example:  $\leq$  on  $\{0, 1\}$ .)

A function  $f : X^k \rightarrow X$  **preserves**  $\rho$  iff:

for all  $\begin{pmatrix} x_1 \\ y_1 \end{pmatrix}, \dots, \begin{pmatrix} x_k \\ y_k \end{pmatrix} \in \rho$ , we have  $\begin{pmatrix} f(x_1, \dots, x_k) \\ f(y_1, \dots, y_k) \end{pmatrix} \in \rho$ .

## Lemma

*If all  $f \in S \subseteq \mathcal{O}_X$  preserve  $\rho$ , then all  $f \in \langle S \rangle$  preserve  $\rho$ .*



# Polymorphisms, definition

## Definition

Let  $\rho \subseteq X^m$  be an  $m$ -ary relation, and let  $f : X^k \rightarrow X$  be a  $k$ -ary function. We say that “ $f$  preserves  $\rho$ ” ( $f \triangleright \rho$ ,  $f \in \text{Pol}(\rho)$ ) if:

- for all  $(a_{i,j} : i \leq m, j \leq k) \in X^{m \times k}$ :

whenever  $a_{*,1} \in \rho, \dots, a_{*,k} \in \rho$ , then also  $\begin{pmatrix} f(a_{1,*}) \\ \vdots \\ f(a_{m,*}) \end{pmatrix} \in \rho$ .

(We let  $a_{*,j} := \begin{pmatrix} a_{1,j} \\ \vdots \\ a_{m,j} \end{pmatrix}$ , similarly  $a_{i,*} = (a_{i,1}, \dots, a_{i,k})$ .)

Equivalently: Letting  $(X^m, f^{[m]})$  be the  $m$ -th power of  $(X, f)$ :  
 $f \triangleright \rho$  iff  $\rho$  is a subalgebra of  $(X^m, f^{[m]})$ .

## Polymorphisms, examples

- ▶ Let  $\rho$  be a nontrivial unary relation, i.e.  $\emptyset \subsetneq \rho \subsetneq X$ . Then  $\text{Pol}(\rho)$  is the set of all operations  $f$  such that  $\rho$  is a **subalgebra** of  $(X, f)$ .
- ▶ Let  $\rho \subseteq X \times X$  be an equivalence relation. Then  $\text{Pol}(\rho)$  is the set of all operations  $f$  such that  $\rho$  is a **congruence relation** of the algebra  $(X, f)$ .
- ▶ Let  $\rho \subseteq X \times X$  be a (reflexive) partial order. Then  $\text{Pol}(\rho)$  is the set of all **pointwise monotone** operations.
- ▶ Let  $\rho \subseteq X \times X$  be the graph of a function  $r$ :  
 $\rho = \{(x, r(x)) : x \in X\}$ .  
Then  $\text{Pol}(\rho)$  is the set of all functions  $f$  such that  $r$  is an **endomorphism** of  $(X, f)$ , i.e.,  $f$  commutes with  $r$ .

Fix a finite base set  $X$ .

### Definition

For any relation  $\rho \subseteq X^m$  let  $\text{Pol}(\rho)$  be the set of all operations preserving  $\rho$ :  $\text{Pol}(\rho) := \{f \in \mathcal{O}_X \mid f \triangleright \rho\}$

For a set  $R$  of relations, let  $\text{POL}(R) := \bigcap_{\rho \in R} \text{Pol}(\rho)$ .

### Lemma

*If  $S \subseteq \text{Pol}(\rho)$ , then also  $\langle S \rangle \subseteq \text{Pol}(\rho)$ . In particular,  $\text{Pol}(\rho)$  and also  $\text{POL}(R)$  are always clones.*

### Theorem

*For every clone  $C \subseteq \mathcal{O}_X$  there exists:*

- ▶ *A set  $S \subseteq \mathcal{O}_X$  such that  $C = \langle S \rangle$ . (Trivial)*
- ▶ *A set  $R$  of relations such that  $C = \text{POL}(R)$ .*

(Helpful to show incompleteness.)

# Galois connection

## Theorem

*For every clone  $C \subseteq \mathcal{O}_X$  there exists a set  $R$  of relations such that  $C = \text{POL}(R) = \{f \mid \forall \rho \in R : f \triangleright \rho\}$ .*

## Proof sketch.

The largest set  $R$  satisfying  $\forall \rho \in R : C \subseteq \text{Pol}(\rho)$  is the set

$$\text{INV}(C) := \{\rho \mid \forall f \in C : f \triangleright \rho\}$$

For finite sets  $X$ , we can check that  $C = \text{POL}(\text{INV}(C))$ .

even:  $\langle S \rangle = \text{POL}(\text{INV}(S))$  for all  $S \subseteq \mathcal{O}_X$ .

We will see a construction of a “better” set  $R$  with  $C = \text{POL}(R)$  later.

## Pol: completeness criterion

Fix a finite base set  $X$ .

### Theorem

*For every clone  $C \subseteq \mathcal{O}_X$  there exists a set  $R$  of relations such that  $C = \text{POL}(R)$ .*

### Corollary

*If  $S \subseteq \mathcal{O}_X$  is not complete (i.e.,  $\langle S \rangle \neq \mathcal{O}_X$ ), then there is a nontrivial relation  $\rho$  such that  $S \subseteq \text{Pol}(\rho)$ , hence  $\langle S \rangle \subseteq \text{Pol}(\rho)$ .*

(But there are so many candidates for  $\rho$ ! Want to search a small set.  $\rightarrow$  precomplete clones)

# Precomplete clones

## Definition

A clone  $C \subseteq \mathcal{O}_X$  is “precomplete” (or “maximal”) if  $C \neq \mathcal{O}_X$ , but there is no clone  $D$  satisfying  $C \subsetneq D \subsetneq \mathcal{O}_X$ .

## Theorem

*For any clone  $C \subsetneq \mathcal{O}_X$  there is a precomplete clone  $C'$  with  $C \subseteq C'$ .*

(Remark: Not true for infinite sets!)

## Proof.

(Use Zorn's lemma??) Let  $\mathcal{O}_X = \langle f \rangle$ . Among all clones  $D$  with  $C \subseteq D$ ,  $f \notin D$ , find a maximal element.

(Better proof: later)

## Examples of precomplete clones

### Example

Let  $\emptyset \subsetneq \rho \subsetneq X$ . Then  $\text{Pol}(\rho)$  is precomplete.

### Proof.

Assuming  $g \notin \text{Pol}(\rho)$ , we let  $C := \langle \text{Pol}(\rho) \cup \{g\} \rangle$ ; we show  $C = \mathcal{O}_X$ .

First show that there is  $b \notin \rho$  such that the constant operation  $c_b$  with value  $b$  is in  $C$ .

For any function  $f : X^k \rightarrow X$  let  $\hat{f} : X^{k+1} \rightarrow X$  be defined by  $\hat{f}(\vec{x}, b) = f(\vec{x})$ , and  $\hat{f}(\vec{x}, y) \in \rho$  arbitrary for  $y \neq b$ . Then  $\hat{f} \in C$ , and  $f(\vec{x}) = \hat{f}(\vec{x}, c_b(x_1))$ , so  $f \in C$ .

### Example

Let  $\rho$  be a bounded partial order. Then  $\text{Pol}(\rho)$  is precomplete.

# Rosenberg's list

## Theorem

*Let  $X = \{1, \dots, k\}$ . Then there is an explicit finite list of relations  $\rho_1, \dots, \rho_m$  (including, for example, all nontrivial unary relations, all bounded partial orders) such that every precomplete clone on  $X$  is one of  $\text{Pol}(\rho_1), \dots, \text{Pol}(\rho_m)$ .*

**Completeness criterion** If  $\langle S \rangle \neq \mathcal{O}_X$  iff there is some  $i$  with  $\forall f \in S : f \triangleright \rho_i$ .



## $k$ -ary fragments

Let  $D$  be a  $k$ -ary clone. The **smallest** clone  $C$  with  $C \cap \mathcal{O}_X^{(k)} = D$  is  $\langle D \rangle$ .

$D \subseteq X^{X^k}$  can be viewed as a relation on  $X$ .

The **largest** clone  $C$  with  $C \cap \mathcal{O}_X^{(k)} = D$  is

$$\text{Pol}(D) = \bigcup_n \{f \in \mathcal{O}_X^{(n)} \mid \forall d_1, \dots, d_n \in D : f(d_1, \dots, d_n) \in D\}$$

For any clone  $E$ , the clones  $\text{Pol}(E \cap \mathcal{O}_X^{(k)})$  approximate  $E$  from above, agreeing with  $E$  on larger and larger sets:

$$\text{Pol}(E \cap \mathcal{O}_X^{(k)}) \cap \mathcal{O}_X^{(k)} = E \cap \mathcal{O}_X^{(k)}.$$

### Theorem

For all clones  $E$ :  $E = \bigcap_k \text{Pol}(E \cap \mathcal{O}_X^{(k)})$ .

## $Cl(X)$ is dually atomic

### Theorem

*Let  $X$  be finite,  $C \neq \mathcal{O}_X$  a clone.*

*Then there is a precomplete clone  $D \supseteq C$ .*

### Proof.

Let  $C' \supseteq C$  be such that  $C' \cap \mathcal{O}_X^{(2)}$  is maximal. (finite!)

Let  $D := \text{Pol}(C')$ .