Ulam Style Cut and Choose Games

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Figure: A piece of cake

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Assume there's a piece of cake, and we want Ann and Bob to fairly share it. An easy way to do so is to let Ann *cut* the cake in two, then let Bob *choose* his piece, and let Ann have the remaining piece.

This principle is already mentioned in the bible, some hundred years BC. Modern investigation of *fair division* was initiated by Steinhaus, Banach and Knaster in the 1940ies. Initally, they extended fair division to a larger number of people, say now you have Ann, Bob and Chris...

In set theory, we tend to be very interested in *infinite games*. Mycielski and Ulam, in the 1960ies, proposed various *infinite cut and choose games*.

We have two players, going by the names of *Cut* and *Choose*. Instead of a piece of cake, we have a set X over which our game takes place. For a start, you may perhaps think of X being either \mathbb{N} or \mathbb{R} .

- In their first move, *Cut* partitions X into two (disjoint) pieces, and *Choose* picks one of them, call it X_0 .
- Now *Cut* partitions X_0 into two pieces, and *Choose* picks one of them, call it X_1 .
- Now Cut partitions X₁ ...

In the end, *Choose* wins a run of the game in case $\bigcap X_i$ has (at least) two distinct elements.

Let us denote the above game as $\mathcal{U}(X)$.

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Why *two* elements?

Assume we declare *Choose* to win if only $\bigcap X_i$ were nonempty. Let X be any set, and let a be an element of X. Now in a run of $\mathcal{U}(X)$ with this modified winning condition, whatever partition *Cut* presents, let *Choose* simply pick the part X_i that contains a as an element. In the end, clearly $a \in \bigcap X_i$, and thus *Choose* would easily win all the time.

This is what we call a *winning strategy*, in this case for *Choose*.

A *strategy* (for one particular player) is a function which receives as input a play of the game so far, and its output is the next move to be made by that player. A strategy is *winning* if following it guarantees a win, regardless of the other player's moves.

In set theory, we like to say that some player *wins* a certain game to express that they have a winning strategy. We say that they win *a particular run* of a game otherwise.

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It is very easy to see that *Cut* wins $\mathcal{U}(\mathbb{N})$ – remember that in a run of the game, players cut and choose infinitely often, and *Choose* wins if the intersection of their choices has **two** distinct elements.

Cut wins by removing one natural number in each step:

- In their first move, they partition $\mathbb N$ into $\{0\}$ and $[1,\infty).$
- Choose will have to pick $[1,\infty)$.
- In their second move, they present the choices $\{1\}$ and $[2,\infty).$
- Choose will have to pick $[2,\infty)$.
- Next options will be $\{2\}$ and $[3,\infty),$...

Obviously, the intersection of all choices will be empty, hence *Cut* wins.

Note that this had nothing to do with any properties of the natural numbers other than their countability!

Cut wins over \mathbb{R}

It is still fairly easy for *Cut* to win $\mathcal{U}(\mathbb{R})$. For notational simplicity, let's play over [0, 1). *Cut* wins by simply cutting in half in each step:

- In their first move, they partition [0,1) into $[0,\frac{1}{2})$ and $[\frac{1}{2},1)$.
- Choose picks one of these intervals [a, b).
- Given [a, b), let *Cut* partition it into $[a, \frac{a+b}{2})$ and $[\frac{a+b}{2}, b)$.
- Choose again picks one of these intervals.

The lengths of these intervals *converge* to 0, so their intersection will contain at most one element. But this means that *Cut* wins.

Note again that this had nothing to do with any particular properties of [0,1) other than its size. If any set stands in bijection with [0,1), like \mathbb{R} for example, then we can still make use of the above argument modulo this bijection. So the only thing that matters here is the *size*, or *cardinality*, of the underlying set. For any set X, let |X| denote its cardinality.

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If $|X| \leq |\mathbb{R}|$, that is there is an injection from X into \mathbb{R} , and thus into [0, 1), then by the above argument, *Cut* has a winning strategy in $\mathcal{U}(X)$. (This means of course that *Choose* cannot have a winning strategy!)

If *Choose* wants to have a chance of winning, they need to play over a reasonably large set, at least larger than $|\mathbb{R}|$.

The question whether *Choose* can ever win is in fact strongly related to that of the possible existence of very large infinities, namely a certain type of *large cardinal*: the existence of *measurable cardinals* (introduced by Ulam in 1930).

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Inaccessible cardinals (Hausdorff 1908, Sierpinski-Tarski-Zermelo 1930) can be seen as the *smallest* of large cardinals. They are sets X of uncountable size |X| with the properties that

- If |Y| < |X|, then also $|\mathcal{P}(Y)| < |X|$, and
- X cannot be written as the union of less than |X|-many sets of size less than X.

Note that $|\mathbb{R}|$ is not inaccessible, nor is $|\mathcal{P}(\mathbb{R})|$, $|\mathcal{P}(\mathcal{P}(\mathbb{R}))|$, ... But also $|\bigcup \mathcal{P}^i(\mathbb{R})|$ is not inaccessible, for it clearly fails to obey the second condition above. In fact, inaccessible cardinals cannot be shown to actually exist. But their existence is often a useful extra assumption. For example, it is equivalent to the existence of *Grothendieck universes*, as used in algebraic geometry.

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Measurable cardinals are uncountable sets κ for which there exists an object (called a *measurable ultrafilter*) deciding for each $A \subseteq \kappa$ whether A or its complement is considered *large* (and the other one *small*), in such a way that the intersection of less than $|\kappa|$ -many large sets will still be large.

Choose wins $\mathcal{U}(\kappa)$ in case κ is a measurable cardinal!

Measurable cardinals are fairly large inaccessible cardinals. The smallest inaccessible cardinal is never measurable. In fact, below any measurable cardinal κ , there are $|\kappa|$ -many smaller inaccessible cardinals.

Theorem (Silver-Solovay, 1970ies)

If there exists a set X such that Choose has a winning strategy in the game $\mathcal{U}(X)$, then there is a generically measurable cardinal, and hence a universe of mathematics in which there exists a measurable cardinal.

Corollary

There are universes of mathematics in which *Choose* has no winning strategy for $\mathcal{U}(X)$ over any set X.

Theorem (Laver, 1970ies)

If a measurable cardinal exists, then there is a universe of mathematics in which Choose has a winning strategy in the game $\mathcal{U}(\mathcal{P}(\mathbb{R}))$, that is the Ulam game on the powerset of \mathbb{R} .

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Leastness is interesting

If Choose has a winning strategy in the game $\mathcal{U}(X)$ and $|X| \leq |Y|$, then Choose has a winning strategy in $\mathcal{U}(Y)$ as well – simply use the strategy for the game $\mathcal{U}(X)$ modulo an injection from X to Y.



So what is interesting is the *smallest* size of a set X so that *Choose* can win $\mathcal{U}(X)$, if there is any at all.

A combination of known results

If there is a measurable cardinal, then there is a universe of mathematics with a measurable cardinal κ , such that *Choose* does not have a winning strategy in $\mathcal{U}(X)$ whenever $|X| < |\kappa|$.

Remember that Choose always has a winning strategy for $\mathcal{U}(\kappa)$ for measurable κ .

One of the main results of our paper is the following:

Theorem

If a measurable cardinal exists, then there is a universe of mathematics in which λ is an inaccessible cardinal that is not measurable (in fact, not even weakly compact), however λ is the least cardinal such that Choose has a winning strategy in the game $\mathcal{U}(\lambda)$.

This means that *Choose* can also *first win* the Ulam game over *fairly small* inaccessible cardinals.

Question

Can *Choose* also first win the Ulam game over the *least* inaccessible cardinal?

Let's go back to the definition of the Ulam game. It seems somewhat annoying having to require two elements in the final intersection in order for *Choose* to win. This is taken care of by the *weak Ulam game* $\overline{U}(X)$. It proceeds just like the Ulam game, however with two modifications:

- In each of their moves, *Choose* has to pick a set with at least two elements, however
- for *Choose* to win, the intersection of all choices only needs to be nonempty.

Observe that simply fixing some $a \in X$ in advance and always picking the part that contains a as an element is not a valid strategy for *Choose* anymore in the game $\overline{\mathcal{U}}(X)$, for it cannot be applied when presented a partition of the form $\langle \{a\}, Y \rangle$. So it doesn't seem that *Choose* can trivially win $\overline{\mathcal{U}}(X)$.

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Remember that Cut had an easy strategy to win $\mathcal{U}([0,1))$, and thus also $\mathcal{U}(\mathbb{R})$ by always cutting given intervals in half. However:

Proposition

Cut does *not* have a winning strategy in $\overline{\mathcal{U}}(\mathbb{R})$.

Proof: By the picture on the next slide...

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Cut does not win $\overline{\mathcal{U}}(\mathbb{R})$



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However things aren't completely different...

Proposition

Choose does not have a winning strategy in the game $\overline{\mathcal{U}}(\mathbb{R})$.

Thus, *neither* player has a winning strategy in the game $\overline{\mathcal{U}}(\mathbb{R})$. We call such games *undetermined*.

In analogy to the classical theorems on $\mathcal{U}(X)$, we can show:

Observation

If there exists a set X such that *Choose* has a winning strategy in the game $\overline{\mathcal{U}}(X)$, then there is a *generically measurable cardinal*, and hence a universe of mathematics in which there exists a measurable cardinal.

Choose does not have a winning strategy in $\bar{\mathcal{U}}(\mathbb{R})$

Proof-Sketch: Assume that *Choose* has a winning strategy σ in the game $\overline{\mathcal{U}}(X)$ and $X \subseteq \mathbb{R}$ with $|X| = |\mathbb{R}|$.

Key technical lemma

 σ is *dynamic*, that is, whenever *Choose* picks some set Y in a play of $\overline{\mathcal{U}}(\mathbb{R})$ according to σ , then there are two distinct (finite) continuations of that play (with *Cut* playing differently) in which *Choose* plays according to σ , such that in the respective last move, *Choose* picks a set Y_0 in one run, and a set Y_1 in another, so that Y_0 and Y_1 are disjoint.

- \Rightarrow X contains a continuous image of the Cantor set.
 - Pick X ⊆ ℝ of size |ℝ| which doesn't contain a continuous image of the Cantor set (exists by an easy diagonalization argument).
 - This means that *Choose* cannot have a winning strategy for $\overline{\mathcal{U}}(X)$.
 - But having a winning strategy or not only depends on the cardinality of the starting set.
 - Therefore, Choose does not have a winning strategy for $ar{\mathcal{U}}(\mathbb{R})$.