

GRAZ, July 2006

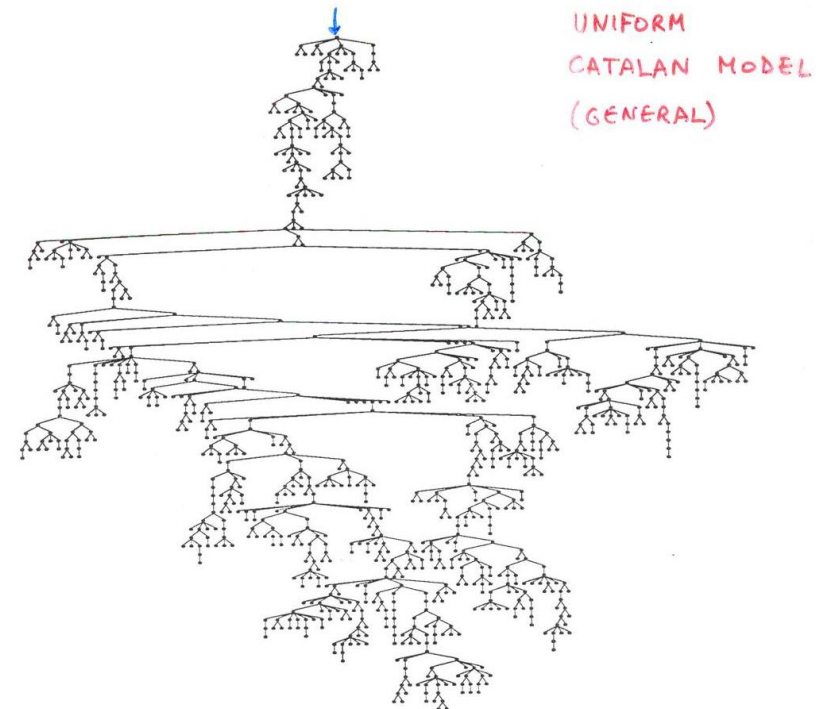
Singularities and Random Combinatorial Structures

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ANALYTIC COMBINATORICS

- Symbolic methods B1
- Complex asymptotic methods B2-B3
- Random structures B4.

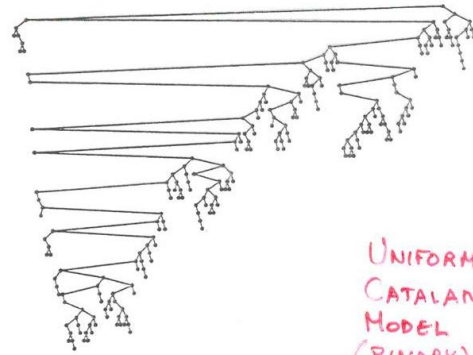
Book by P.F + R. Sedgewick (2007)



§5.4

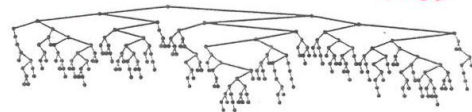
TREES

231



UNIFORM CATALAN MODEL (BINARY)

Figure 5.5 A random binary tree with 256 internal nodes



BINARY SEARCH TREE

Figure 5.11 A binary search tree built from 256 randomly ordered keys

Sedgewick & Flajolet: An Introduction to the Analysis of Algorithms

1



1, 1, 2, 5, 14, ...

1



?

2



5



14

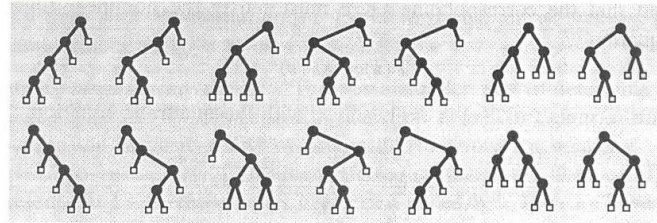
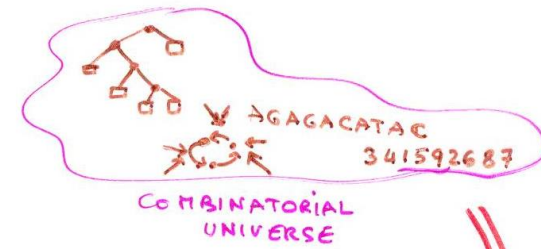
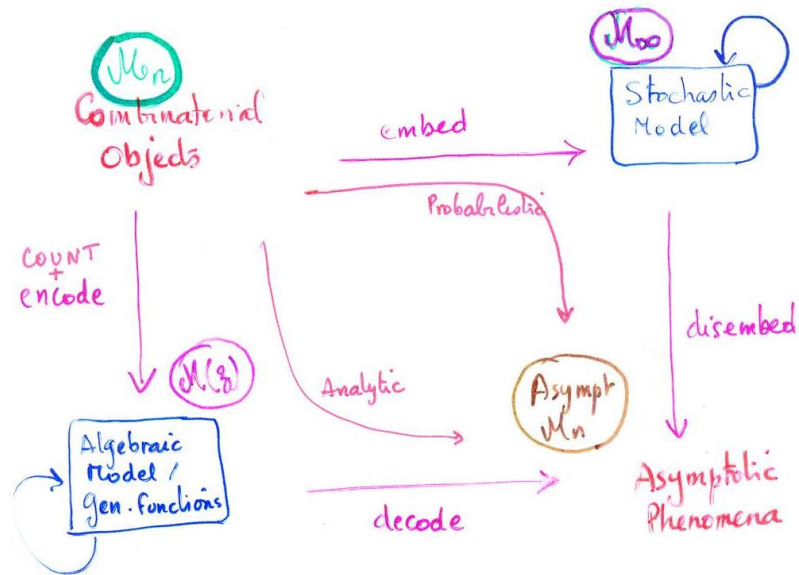


Figure 3.1 All binary trees with 1, 2, 3, 4, and 5 external nodes

PART I. SYMBOLIC METHODS



Algebra of Generating Functions:

$$\sum_{n=0}^{\infty} \frac{1}{n!} \binom{2n}{n} z^n = \frac{1-\sqrt{1-4z}}{2z}$$

$$\sum_{n=0}^{\infty} 2^n z^n = \frac{1}{1-2z}$$

$$\sum_{n=0}^{\infty} n! z^n = e^z \dots$$

ALGEBRA OF GENERATING FUNCTIONS

CHAPTER 1Unlabelled Structures and Ordinary
Generating Functions

Ordinary Generating Function OGF

$$(f_n) \longrightarrow f(z) = \sum_{n=0}^{\infty} f_n z^n$$

number sequence
e.g., counting sequence

(Later: Exponential Generating Function (EGF))

$$(f_n) \longrightarrow \sum f_n \frac{z^n}{n!}$$

\mathcal{C} = a combinatorial class (at most denumerable set
w/ size function)

\mathcal{C}_n = subclass of objects of size n

C_n = # objects of size n = card(\mathcal{C}_n)

$$C(z) = \sum_{n=0}^{\infty} C_n z^n = \sum_{\gamma \in \mathcal{C}} z^{|\gamma|}$$

Naturally: Count up to combinatorial isomorphism

($\mathcal{C} \equiv \mathcal{D}$ if \exists size-preserving bijection).

How many binary trees

B_N with
 N external nodes?

$$B = \square + \begin{array}{c} \bullet \\ / \quad \backslash \\ B \quad B \end{array}$$

type BinTree = Union(
~~External Node~~,
Product (BinTree, BinTree));

Classical (naïve?) approach

$$B_N = \sum_{k=1}^{N-1} B_k B_{N-k} \quad (N > 1)$$

$$\begin{aligned} B(z) &:= \sum_N B_N z^N \\ &= \sum_{t \in B} z^{|t|} \end{aligned}$$

$$\begin{aligned} \sum_{t \in B} z^{|t|} &= \sum_{t = \square} z^{|t|} + \sum_{t_1, t_2} z^{|t_1| + |t_2|} \\ &= z^{|\square|} + \sum_{t_1 \in B} z^{|t_1|} \times \sum_{t_2 \in B} z^{|t_2|} \end{aligned}$$

$$B(z) = z + (B(z) \times B(z))$$

$$B = z + B^2$$

Solve the quadratic equation

$$B(z) = \frac{1 - \sqrt{1 - 4z}}{2} = \frac{1}{2} - \frac{1}{2}(1 - 4z)^{1/2}$$

The number of binary trees with N external nodes is

$$B_N = \frac{1}{N} \binom{2N-2}{N-1} \quad (\text{Catalan numbers})$$

OUTLINE

Define a collection of *sub-theoretic combinatorial*
CONSTRUCTIONS

union, product, sequence, set, cycle (...)

allowing RECURSIVE DEFINITIONS

Then:

meta-THM1: Generating functions are
automatically computable (equations!)

meta-THM2: Counting sequences are also
automatically computable $O(N^2)$
 $\hookrightarrow O(N \log N)$

meta-THM3: Random generation is
fast $O(N \log N)$
arithmetic complexity

Theorem 1.1 There exists a dictionary

{ Combinatorial
Constructions } \leftrightarrow { OGF
operations }

UNLABELLED

Construction	Translation (OGF)
$\mathcal{F} = \mathcal{G} \cup \mathcal{H}$	$F(z) = G(z) + H(z)$
$\mathcal{F} = \mathcal{G} \times \mathcal{H}$	$F(z) = G(z) \cdot H(z)$
$\mathcal{F} = \text{sequence}(\mathcal{G}) \equiv \mathcal{G}^*$	$F(z) = \frac{1}{1-G(z)}$
$\mathcal{F} = \text{set}(\mathcal{G})$	$F(z) = \exp(G(z) - \frac{1}{2}G(z)^2 + \frac{1}{3}G(z)^3 - \dots)$
$\mathcal{F} = \text{multiset}(\mathcal{G})$	$F(z) = \exp(G(z) + \frac{1}{2}G(z)^2 + \frac{1}{3}G(z)^3 + \dots)$
$\mathcal{F} = \text{cycle}(\mathcal{G})$	$F(z) = \log(1 - G(z))^{-1} + \dots$
$\mathcal{F} = \mathcal{G}[\mathcal{H}]$	$F(z) = G(H(z))$

\mathcal{E} or $\mathbf{1}$: neutral class formed with one element of
size 0 [cf empty word] $\rightarrow E(z) \equiv 1$.

\mathcal{Z} : atomic class formed with one element of
size 1, an ATOM. $\rightarrow Z(z) \equiv z$

Proof (s)

$$A \longmapsto A(z) = \sum A_n z^n$$

$$A(z) = \sum_{\alpha \in A} z^{|\alpha|}$$

UNION

$$C = A \cup B$$

$$C(z) = A(z) + B(z)$$

$$\sum_{\gamma \in C} z^{|\gamma|} = \sum_{\alpha \in A} z^{|\alpha|} + \sum_{\beta \in B} z^{|\beta|}$$

CARTESIAN PRODUCT

$$C = A \times B$$

$$C(z) = A(z) \times B(z)$$

$$\sum_{\gamma \in C} z^{|\gamma|} = \sum_{(\alpha, \beta) \in A \times B} z^{|\alpha| + |\beta|}$$

SEQUENCE

$$C = \text{Seq}(A)$$

$$C(z) = \frac{1}{1 - A(z)}$$

$$(i) C = 1 + (A \times C) \quad (ii) \frac{1}{1-f} = 1 + f + f^2 + \dots$$

SET

$$C = \text{Set}(A)$$

$$\cong \prod_{\alpha \in A} (1 + f^\alpha)$$

$$C(z) = \exp(A(z) - \frac{A(z)^2}{2} + \dots)$$

Pólya Operator

$$C(z) = \prod_{n=1}^{\infty} (1 + z^n)^{A_n} = \exp\left(\sum_n A_n \log(1 + z^n)\right) = \dots$$

MULTISET

$$C = \text{MSet}(A) \cong \prod_{\alpha \in A} \text{Seq}(\{f^\alpha\})$$

$$C(z) = \prod_{n=1}^{\infty} (1 - z^n)^{-A_n} = \exp(A(z) + \frac{1}{2} A(z)^2 + \dots)$$

End of Proof ■

This theorem permits us to write **AUTOMATICALLY** for binary trees

$$B = \blacksquare + (B \circlearrowleft B)$$



$$B(z) = z + B(z)^2$$

EXAMPLE 1 WORDS

$$W = \text{Seq}(\underline{a} + \underline{b})$$

$$W(z) = \frac{1}{1-2z}$$

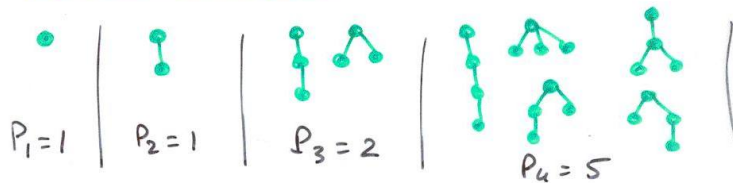
Starting with \underline{b} , never more than 4 consec. " \underline{a} "

$$W = \text{Seq}(\underline{b} \times (1 + \underline{a} + \underline{a}\underline{a} + \underline{a}\underline{a}\underline{a} + \underline{a}\underline{a}\underline{a}\underline{a}))$$

$$W(z) = \frac{1}{1-(z+z^2+z^3+z^4+z^5)}$$

EXAMPLE 2 PLANE TREES ("general")

$$P = Z \times \text{Seq}(P)$$



$$P(z) = \frac{z}{1-P(z)} \Rightarrow P(z) = \frac{1-\sqrt{1-4z}}{2} \quad P_n = \frac{1}{n} \binom{2n-2}{n-1}$$

EXAMPLE 3 NONPLANE TREES

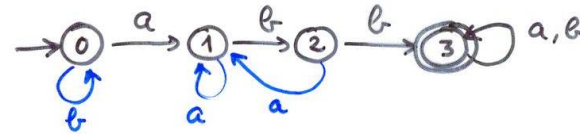
$$U = Z \times \text{Set}(U)$$

$U_1=1, U_2=1, U_3=2, U_4=4 \dots$

$$U(z) = z \cdot \exp\left(\frac{U(z)}{1} + \frac{U(z^2)}{2} + \frac{U(z^3)}{3} + \dots\right)$$

CAYLEY \rightarrow Recurrences; Pólya \rightarrow Asymptotics

Words containing a pattern

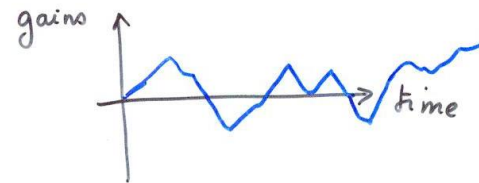


\mathcal{L}_j : language accepted from state j

$$\mathcal{L}_1 = a \cdot \mathcal{L}_1 + b \cdot \mathcal{L}_2$$

Reg \mapsto $\mathbb{Q}(z)$
rational

Walks and Excursions



Walks \mapsto $\mathbb{Q}(z, \sqrt{1-4z^2})$

Excursion =
= $\text{Seq}(\uparrow \text{Excursion} \downarrow)$

Positive path = Excursion \times $\text{Seq}(\uparrow \text{Excursion})$

Draw game = $\text{Seq}(\downarrow \text{Excursion} \uparrow + \uparrow \text{Excursion} \downarrow)$ etc

Exercise: Integer compositions

- a sequence (x_1, \dots, x_k) with $x_j \in \mathbb{Z}_{\geq 1}$ is a composition of n if $x_1 + \dots + x_k = n$. (Order of summands count)

Q: How many compositions of integer n , for $n = 1, 2, 3, 4$
 – for general n ?

Hint: Argue that $\mathcal{C} = \text{Seq}(\mathcal{W})$; $\mathcal{W} = \mathbb{Z} \times \text{Seq}(\mathbb{Z})$

Exercise: Denumerants. In how many ways can you arrange to give a change for n Euro-cents knowing that you have an arbitrary supply of coins of 1¢, 2¢, 5¢, 10¢?

[order of coins does NOT matter!]
 Q1: Prove that OGF is $\frac{1}{(1-z)(1-z^2)(1-z^5)(1-z^{10})}$.

Q2: generalize to Ω . Q3 How to obtain coefficients? (complexity?)

Exercise: Unary-binary trees are (plane) trees where nodes have degree either 0, 1, or 2.

Q1: How many U-B trees of size 1, 2, 3, 4?

$$\text{Using } \begin{cases} \mathcal{U} = \bullet \times (1 + \mathcal{U} + \mathcal{U} \times \mathcal{U}) \\ \mathcal{U} = \bullet + \uparrow \mathcal{U} + \downarrow \mathcal{U} \end{cases}$$

Find equation for OGF and solve it.

Q2: Generalization: let $\Omega \ni 0$ and let \mathcal{Y} be the class of trees with nodes having (out)degrees constrained by Ω . Establish for the OGF the equation

$$(\mathcal{Y} \equiv \mathcal{Y}(z)) \quad \mathcal{Y} = z \phi(\mathcal{Y}).$$

Lagrange Inversion Theorem:

$[z^n] \mathcal{Y}(z) = \frac{1}{n} [w^{n-1}] \phi(w)^n$.
 Q3: Ω finite $\Rightarrow \mathcal{Y}_n$ is a sum of products of ~~binomial~~ multinomial coefficients.

- Exercise:
- Binary trees with n internal nodes $\hat{B}_n = \frac{1}{n+1} \binom{2n}{n}$
 - General (plane) trees with $n+1$ internal nodes $G_{n+1} = \frac{1}{n+1} \binom{2n}{n}$
 - Simple excursions $\uparrow \downarrow$ of length $2n$ $E_{2n} = \frac{1}{n+1} \binom{2n}{n}$
- are equinumerous.

Question why?

[Find a complete set of bijections]

