

# Equations in algebras induced by beautiful first-order structures

**Michael Pinsker**

Technische Universität Wien

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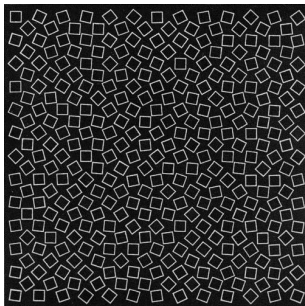
Univerzita Karlova v Praze

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# Outline

- I:** Equations in algebras
- II:** Algebras as invariants of first-order structures
- III:** Compact is the new beautiful?
- IV:** The future



## I: Equations in finite algebras

## Equations / identities

Algebra:  $\mathbf{A} = (A, (f_i)_{i \in I})$ ; each  $f_i: A^{n_i} \rightarrow A$ .

*Equation / identity:*

Formal expression  $s \approx t$ , for abstract *terms*  $s, t$  over some functional language  $\tau$ , e.g.

$$g(g(x, y), z) \approx g(x, g(y, z)) .$$

Equation *is satisfied* in  $\mathbf{A}$  if  $g$  can be assigned a *term function*  $\mathbf{g}$  over  $\mathbf{A}$  such that

$$\mathbf{A} \models \forall x \forall y \forall z \mathbf{g}(g(x, y), z) = \mathbf{g}(x, g(y, z)) .$$

Independent of language of  $\mathbf{A}$ .

### Examples.

- The above equation satisfied in any  $\mathbf{A}$  by a projection.
- The equations  $s(y, y, x) = s(x, y, y) = x$  are satisfied in any group  $(A, +, -)$ . [Set  $\mathbf{s}(x, y, z) := x - y + z$ ].

# Clones

**Universal Algebra:** Equations satisfied in  $\mathbf{A} \iff$  structure of  $\mathbf{A}$ .

**Examples:** Congruences of  $\mathbf{A}$ .

- $\mathbf{A}$  satisfies  $s(x, y, y) \approx s(y, y, x) \approx x \implies$   
 $\mathbf{A}$  has permuting congruences (CP)(+ all algebras in *variety* of  $\mathbf{A}$ ).  
Example: groups  $(A, +, -)$ . (*Mal'cev '54*)
- $\mathbf{A}$  satisfies *near unanimity* (*nu*) equations

$$n(x, \dots, x, y) \approx n(x, \dots, x, y, x) \approx \dots \approx n(y, x, \dots, x) = x$$

$\implies \mathbf{A}$  (+ its variety) is congruence distributive (CD).

Example: lattices. Equivalent to CD: *Jónsson '68 equations*.

*Term clone*  $\text{Clo}(\mathbf{A})$  of algebra  $\mathbf{A}$ : smallest set of functions on  $A$

- containing fundamental operations, projections
- closed under composition.

**Goal:** Structure of  $\mathbf{A} \iff$  composition structure of  $\text{Clo}(\mathbf{A})$ .

# Clone homomorphisms

$\text{Clo}(\mathbf{A}) \subseteq \text{Clo}(\mathbf{B}) \implies \mathbf{B}$  has more structure.

More generally:

## Definition

$\xi: \text{Clo}(\mathbf{A}) \rightarrow \text{Clo}(\mathbf{B})$  is a *clone homomorphism* if it preserves

- arities;
- projections;
- composition.

*Quasiorder* on algebras. The larger  $\mathbf{A}$ , the more structure.

**A smallest element:** *clone*  $\mathbf{0}$  of projections on set of  $\geq 2$  elements.

## Important question

Consequences for  $\mathbf{A}$  when  $\text{Clo}(\mathbf{A}) \not\cong \mathbf{0}$ ? Draw picture. Say why.

# Non-triviality + Idempotency

## Definition

$\mathbf{A}$  is *idempotent* if  $t(x, \dots, x) \approx x$  holds for all  $t \in \text{Clo}(\mathbf{A})$ .

**Why idempotent?** Practical + Philosophical reasons + Evidence.

## Theorem (Olšák '16)

Let  $\mathbf{A}$  be idempotent, and  $\text{Clo}(\mathbf{A}) \not\approx \mathbf{0}$ . Then  $\mathbf{A}$  satisfies

$$t(xyy, yxx) \approx t(yxy, xyx) \approx t(yyx, xxy).$$

## Theorem (Maróti + McKenzie '08)

Let  $\mathbf{A}$  be idempotent, in *locally finite variety*,  $\text{Clo}(\mathbf{A}) \not\approx \mathbf{0}$ .

Then  $\mathbf{A}$  satisfies for some arity  $n \geq 2$

$$w(x, \dots, x, y) \approx w(x, \dots, x, y, x) \approx \dots \approx w(y, x, \dots, x).$$

(*weak near unanimity term*)

## Non-triviality + Idempotency + finiteness

Theorem (Siggers '11; Kearnes + Marković + McKenzie '14)

Let  $\mathbf{A}$  be idempotent, in locally finite variety, and  $\text{Clo}(\mathbf{A}) \not\cong \mathbf{0}$ .  
Then  $\mathbf{A}$  satisfies

$$s(x, y, x, z, y, z) \approx s(y, x, z, x, z, y) .$$

and

$$q(a, r, e, a) \approx q(r, a, r, e) .$$

(6-ary and 4-ary Siggers terms)

Theorem (Barto + Kozik '11)

Let  $\mathbf{A}$  be idempotent, *finite*, and  $\text{Clo}(\mathbf{A}) \not\cong \mathbf{0}$ . Then  $\mathbf{A}$  satisfies

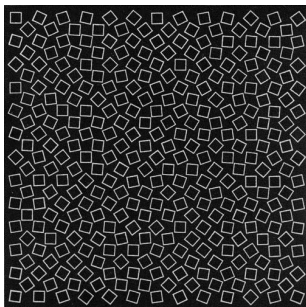
$$c(x_1, \dots, x_n) \approx c(x_2, \dots, x_n, x_1) \quad \text{for some } n \geq 2 .$$

(Cyclic term)



# Summary of Part I

- Equations of  $\text{Clo}(\mathbf{A}) \Leftrightarrow$  Structure of  $\mathbf{A}$  (e.g. congruences)
- Idempotency + **any** equations ( $\text{Clo}(\mathbf{A}) \not\cong \mathbf{0}$ )  $\Rightarrow$  Structure
- Add finiteness conditions  $\Rightarrow$  More structure



## II: Algebras from first-order structures

# Clones as symmetries

Let  $\mathbb{A} = (A, (R_i)_{i \in I})$  be a first-order structure with relations  $R_i$  (wlog).

## Symmetries:

- $\text{Aut}(\mathbb{A}) \subseteq \text{Sym}(A)$  ... automorphism group of  $\mathbb{A}$ ;
- $\text{End}(\mathbb{A}) \subseteq A^A$  ... endomorphism monoid of  $\mathbb{A}$ ;
- $\text{Pol}(\mathbb{A}) \subseteq \bigcup_{n \geq 1} A^{A^n}$  ... polymorphism clone of  $\mathbb{A}$ .

$$\text{Pol}(\mathbb{A}) := \{f: A^n \rightarrow A \mid f \text{ homomorphism}\}.$$

## Topology:

$A$  ... discrete,  $A^{A^n}$  ... product topology (pointwise convergence),  
 $\bigcup_{n \geq 1} A^{A^n}$  ... sum space.

A permutation group / transformation monoid / term clone  
is of the form  $\text{Aut}(\mathbb{A}) / \text{End}(\mathbb{A}) / \text{Pol}(\mathbb{A}) \Leftrightarrow$  it is *topologically closed*.

# Galois-correspondence

The assignments

$$\mathbb{A} \mapsto \text{Pol}(\mathbb{A})$$

$$\mathbf{A} \mapsto \text{Inv}(\mathbf{A})$$

define a Galois-correspondence ( $\text{Inv}(\mathbf{A})$ ... invariant relations).

First-order structures  $\leftrightarrow$  algebras.

**Roughly:**

- $\text{Aut}(\mathbb{A})$ ... first-order structure of  $\mathbb{A}$ .
- $\text{Pol}(\mathbb{A})$ ... primitive positive structure of  $\mathbb{A}$  (finer).

More conditions of beauty for  $\mathbf{A} = \text{Pol}(\mathbb{A})$  via correspondence:

- Finitely related, i.e.,  $\mathbb{A}$  finite + finite language.
- $\mathbb{A}$   $\omega$ -categorical.
- $\mathbb{A}$  Ramsey.
- $\mathbb{A}$  definable in finitely bounded homogeneous structure...

## Finitely related algebras (classical)

### Conjecture (Valeriote)

Let  $\mathbf{A}$  be in *congruence modular variety* + finitely related.

Then the number of subalgebras of  $\mathbf{A}^n$  grows only exponentially with  $n$ .

### Conjecture (Zádori)

Let  $\mathbf{A}$  be in *congruence distributive variety* + finitely related.

Then  $\mathbf{A}$  satisfies near unanimity equations.

$\text{CSP}(\mathbb{A})$ : computational problem of deciding primitive positive theory of  $\mathbb{A}$ .

For  $\mathbb{A}$  finite: in NP.

### Conjecture (Feder + Vardi '92)

Let  $\mathbb{A} = (\mathbf{A}, R_1, \dots, R_n)$  be finite, with (wlog) idempotent  $\text{Pol}(\mathbb{A})$ .

If  $\text{Pol}(\mathbb{A}) \not\rightarrow \mathbf{0}$ , then  $\text{CSP}(\mathbb{A})$  is polynomial-time solvable.

# Finitely related algebras

## Theorem (Barto '12)

Let  $\mathbf{A}$  be in *congruence modular variety* + finitely related.  
Then the number of subalgebras of  $\mathbf{A}^n$  grows only exponentially with  $n$ .

## Corollary (Barto '10)

Let  $\mathbf{A}$  be in *congruence distributive variety* + finitely related.  
Then  $\mathbf{A}$  satisfies near unanimity equations.

$\text{CSP}(\mathbb{A})$ : computational problem of  
deciding primitive positive theory of  $\mathbb{A}$ .

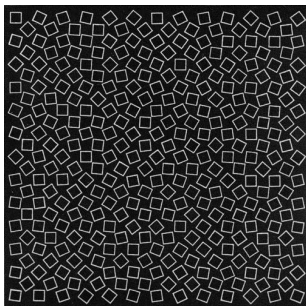
For  $\mathbb{A}$  finite: in NP.

## Theorem (Bulatov '17, Zhuk '17)

Let  $\mathbb{A} = (\mathbf{A}, R_1, \dots, R_n)$  be finite, with (wlog) idempotent  $\text{Pol}(\mathbb{A})$ .  
If  $\text{Pol}(\mathbb{A}) \not\rightarrow \mathbf{0}$ , then  $\text{CSP}(\mathbb{A})$  is polynomial-time solvable.

# Summary of Part II

- Algebras  $\Leftrightarrow$  First-order structures
- Provides more notions of beauty
- Finitely related: works



### III: Compact is the new finite



## Beautiful structures: compactness

Let  $\mathbb{A}$  be countable  $\omega$ -categorical.

$\Leftrightarrow \mathbb{A}^n$  has finitely many first-order definable subsets for every  $n \geq 1$

$\Leftrightarrow \text{Aut}(\mathbb{A}) \curvearrowright \mathbb{A}^n$  has finitely many orbits for every  $n \geq 1$ .

$\Leftrightarrow \text{Aut}(\mathbb{A})$  is *oligomorphic*.

$\mathbb{A}$  “small”, so  $\text{Pol}(\mathbb{A})$  “large” (not locally finite ...). But:

### Observation

For  $f, g \in \text{Pol}(\mathbb{A})$ , set

$$f \sim g : \Leftrightarrow f \in \overline{\{\alpha \circ g \mid \alpha \in \text{Aut}(\mathbb{A})\}}$$

If  $\mathbb{A}$  is  $\omega$ -categorical, then  $\text{Pol}(\mathbb{A})/\sim$  is compact.

Factoring compatible with continuous clone homomorphisms

$\text{Pol}(\mathbb{A}) \rightarrow \text{Pol}(\mathbb{B})$ .

# Topological Birkhoff

$\text{HSP}^{\text{fin}}(\text{Clo}(\mathbf{A}))$ . . . all clones (“actions”) obtained from  $\text{Clo}(\mathbf{A})$  by

- factoring the domain:  $\text{Clo}(\mathbf{A}) \curvearrowright A / \sim$
- restricting the domain:  $\text{Clo}(\mathbf{A}) \curvearrowright A' \subseteq A$
- taking finite powers of the domain:  $\text{Clo}(\mathbf{A}) \curvearrowright A^n$

## Theorem (Birkhoff '35)

Let  $\mathbf{A}$  be finite.

$\text{Clo}(\mathbf{A}) \rightarrow \mathbf{0} \iff \mathbf{0} \in \text{HSP}^{\text{fin}}(\text{Clo}(\mathbf{A})).$

## Theorem (Bodirsky + P. '11)

Let  $\mathbb{A}$  be  $\omega$ -categorical.

$\text{Pol}(\mathbb{A}) \rightarrow \mathbf{0}$  continuously  $\iff \mathbf{0} \in \text{HSP}^{\text{fin}}(\text{Pol}(\mathbb{A})).$

In this situation,  $\mathbb{A}$  is rich:

All finite structures have a *primitive positive interpretation* in  $\mathbb{A}$ .

$\text{CSP}(\mathbb{A})$  is NP-hard.

## Topological non-triviality

Let  $\mathbb{A}$  be  $\omega$ -categorical,  $\text{Pol}(\mathbb{A}) \not\rightarrow \mathbf{0}$  continuously.

Wish to derive equations satisfied in  $\text{Pol}(\mathbb{A})$ .

**Problem:** anti-idempotency of  $\text{Pol}(\mathbb{A})$  implied by  $\omega$ -categoricity.

**Approach I: Approximate idempotency.**

Assume  $\overline{\text{Aut}(\mathbb{A})} = \text{End}(\mathbb{A})$ , consider *stabilizers*  $\text{Pol}(\mathbb{A}, a_0, \dots, a_n)$ .

**Theorem** (Barto + P. '16)

Let  $\mathbb{A}$  be  $\omega$ -categorical,  $\overline{\text{Aut}(\mathbb{A})} = \text{End}(\mathbb{A})$ . TFAE:

- No stabilizer of  $\text{Pol}(\mathbb{A})$  maps continuously to  $\mathbf{0}$
- No stabilizer of  $\text{Pol}(\mathbb{A})$  maps to  $\mathbf{0}$
- $\text{Pol}(\mathbb{A})$  satisfies

$$e \circ s(x, y, x, z, y, z) \approx e' \circ s(y, x, z, x, z, y)$$

**Remark.** Compatible with CSP; interpretations with parameters.

# Linear non-triviality

**Approach I: Approximate idempotency.**

**Approach II: Ignore idempotency.**

Based on observation that most interesting equations are *linear* (non-nested).

Consider *clonoid* homomorphisms  $\text{Clo}(\mathbf{A}) \rightsquigarrow \text{Clo}(\mathbf{B})$  preserving **only linear** equations. Ignore unary functions!

What if  $\text{Pol}(\mathbb{A}) \not\rightsquigarrow \mathbf{0}$ ?

**Theorem** (Barto + Opršal + P. '16; Barto + P. '16; Pham '17)

Let  $\mathbb{A}$  be  $\omega$ -categorical.

If  $\text{Pol}(\mathbb{A}) \rightsquigarrow \mathbf{0}$  unif. continuous, then  $\mathbb{A}$  rich (e.g.  $\text{CSP}(\mathbb{A})$  is NP-hard).

If  $\text{Pol}(\mathbb{A}) \not\rightsquigarrow \mathbf{0}$  unif. continuous, then  $\text{Pol}(\mathbb{A})$  satisfies

$$e \circ s(x, y, x, z, y, z) \approx e' \circ s(y, x, z, x, z, y) .$$

# Orbit growth: extremely beautiful structures

Converse?

Theorem (Olšák + Barto + Kompatscher + Pham + P. '17)

There exists  $\omega$ -categorical  $\mathbb{A}$  with  $\overline{\text{Aut}(\mathbb{A})} = \text{End}(\mathbb{A})$  such that:

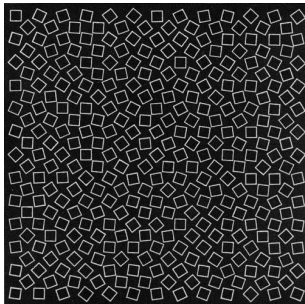
- $\text{Pol}(\mathbb{A})$  satisfies  $e \circ s(x, y, x, z, y, z) \approx e' \circ s(y, x, z, x, z, y)$
- $\text{Pol}(\mathbb{A}) \rightsquigarrow \mathbf{0}$  uniformly continuous.

Theorem (Olšák + Barto + Kompatscher + Pham + P. '17)

The above imply that  $\text{Aut}(\mathbb{A})$  has doubly exponential orbit growth.

Hence, for  $\mathbb{A}$  with slower orbit growth:

$\text{Pol}(\mathbb{A}) \rightsquigarrow \mathbf{0}$  uniformly cont.  $\Leftrightarrow$  no pseudo-Siggers equation.



## IV: The future

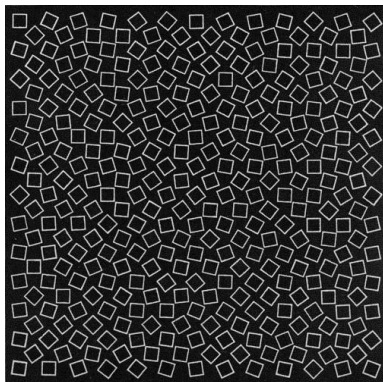
# Summary

## Summary:

- Idempotent / locally finite / finite / finitely related:  
abundance of recent deep “*equations*  $\Leftrightarrow$  *structure*” theorems
- $\omega$ -categoricity yields compactness:  
Work “modulo group” and obtain “*equations*  $\Leftrightarrow$  *structure*”.

## Future:

- Role of topology: which structural properties are algebraic?  
 $\text{Pol}(\mathbb{A}) \rightsquigarrow \mathbf{0} \Leftrightarrow \text{Pol}(\mathbb{A}) \rightsquigarrow \mathbf{0}$  unif. continuously?
- How far does compactness approach go?  
6-ary pseudo-Siggers  $\implies$  4-ary pseudo-Siggers?
- Equations are “fixed points” of clone actions in lack of space.  
General dynamical / Ramsey theoretic principle behind this?
- News: 6-ary Siggers  $\implies$  4-ary Siggers **in general!** (Olšák '18)



Gerhard von Graevenitz     *Regularity - Irregularity V*