

Canonical functions and the Ramsey property, revisited

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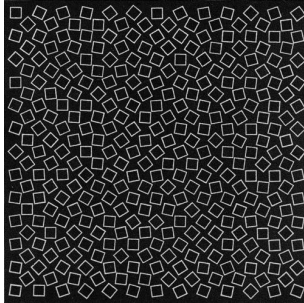
Canonical functions

I: Definition

II: Use

III: How to find

IV: How to avoid



I: What are canonical functions?

Regular behaviour of functions

Example 1

Let $f: A \rightarrow B$ be a function.

There exists an infinite $S \subseteq A$ on which f is injective or constant.

- f behaves “regularly” on S .

Example 2

Let $f: (\mathbb{N}, <) \rightarrow (\mathbb{Q}, <)$. There exists an infinite $S \subseteq \mathbb{N}$ on which f is strictly increasing, decreasing, or constant.

- f behaves on S “regularly” wrt the orders $(\mathbb{N}, <)$ and $(\mathbb{Q}, <)$.
- Note: f does not *preserve* the structure $(\mathbb{N}, <)$ on S .

Example 3

Let $f: (\mathbb{Q}, <, 0) \rightarrow (\mathbb{Q}, <)$. Every finite substructure of $(\mathbb{Q}, <, 0)$ has an isomorphic copy on which f behaves “regularly”.

Canonical functions

Let \mathbb{A}, \mathbb{B} be first-order structures, and $f: \mathbb{A} \rightarrow \mathbb{B}$.

Definition (model-theoretic)

f is **canonical** $:\Leftrightarrow$

f sends tuples of the same type in \mathbb{A} to tuples of the same type in \mathbb{B} .

Definition (algebraic)

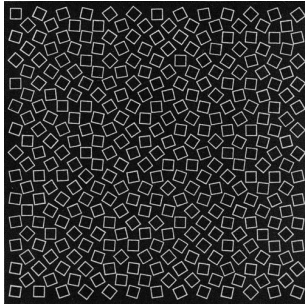
f is **canonical** $:\Leftrightarrow$

\forall tuples \bar{t} in \mathbb{A} $\forall \alpha \in \text{Aut}(\mathbb{A}) \exists \beta \in \text{Aut}(\mathbb{B})$ such that

$$f(\bar{t}) = \beta f \alpha^{-1}(\bar{t}).$$

Definitions equivalent when \mathbb{A}, \mathbb{B} are countable ω -categorical, i.e., $\text{Aut}(\mathbb{A})$ and $\text{Aut}(\mathbb{B})$ act on n -tuples with finitely many orbits:

Tuples have same type \Leftrightarrow they are in the same orbit.



II: The use of canonical functions

Reducts

Theorem (Thomas '91)

Let (V, E) be the random graph.

There exist precisely 3 **reducts of (V, E)** , i.e., closed permutation groups \mathcal{G} with $\text{Aut}(V, E) \subsetneq \mathcal{G} \subsetneq \text{Sym}(V)$:

- $\text{Aut}(V, E) + \gamma \in \text{Sym}(V)$ which flips edges and non-edges.
- $\text{Aut}(V, E) + \delta \in \text{Sym}(V)$ which flips edges and non-edges around a fixed vertex $v \in V$.
- Their join.

Remarks.

- $\gamma: (V, E) \rightarrow (V, E)$ is canonical!
- $\delta: (V, E, v) \rightarrow (V, E)$ is canonical!

Coincidence?

Analyzing reducts

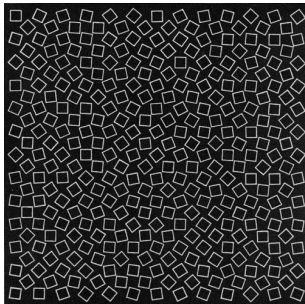
Proof of Thomas' theorem.

- Let $\text{Aut}(V, E) \subsetneq \mathcal{G} \subsetneq \text{Sym}(V)$ be closed.
- Pick $f \in \mathcal{G} \setminus \text{Aut}(V, E)$.
- There is an edge $(a, b) \in V^2$ which f sends to a non-edge.
- Elsewhere f is *ugly*.
- **There seems no way out.**
- But suppose somebody steps in and tells us that $f: (V, E, a, b) \rightarrow (V, E)$ is *canonical*.
- Examine all possible *behaviours* for f .
(Deletes edges around a ? Other edges? Non-edges?...)
- Conclude that in each case, \mathcal{G} contains one of the two groups.
- Continue if \mathcal{G} is even bigger ...

Q.E.D.

More motivation

- Canonicity convenient when one has finite information about a function $f: \mathbb{A} \rightarrow \mathbb{B}$;
want f to be nice outside that information.
- **Polymorphisms** $f: \mathbb{A}^n \rightarrow \mathbb{A}$ for **Constraint Satisfaction Problems**:
Every structure \mathbb{A} defines computational problem:
Given a *primitive positive* sentence ϕ , does $\mathbb{A} \models \phi$ hold?
The more polymorphisms $f: \mathbb{A}^n \rightarrow \mathbb{A}$ has, the easier the problem.
Use polymorphism f to prove polynomial-time algorithm.
Good if f is canonical. . .



III: Canonical functions in an ideal world

Obtaining canonical functions

Let $f: \mathbb{A} \rightarrow \mathbb{B}$.

- Where?
- How?

Definition

The **orbit** of f (under action of $\text{Aut}(\mathbb{B}) \times \text{Aut}(\mathbb{A})$ on $\mathbb{B}^{\mathbb{A}}$):

$$O(f) := \{\beta f \alpha^{-1} \mid \beta \in \text{Aut}(\mathbb{B}), \alpha \in \text{Aut}(\mathbb{A})\}$$

Note: All functions or no function in an orbit are canonical.
But we may be lucky in the (pointwise) closure $\overline{O(f)}$ in $\mathbb{B}^{\mathbb{A}}$.

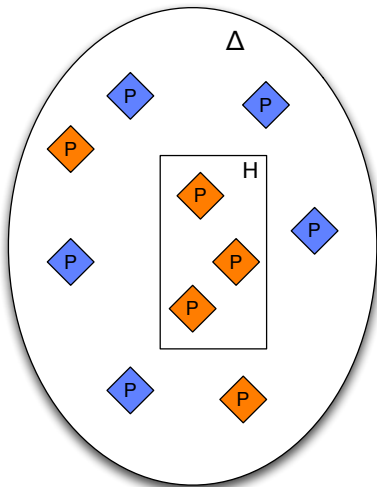
Definition

\mathbb{A} is **Ramsey** $:\Leftrightarrow$

\forall finite $\mathbb{P}, \mathbb{H} \subseteq \mathbb{A} \quad \forall \chi: \binom{\mathbb{A}}{\mathbb{P}} \rightarrow \mathbf{2} \quad \exists \mathbb{H}' \cong \mathbb{H}$ on which χ is constant.

Examples in beginning: Ramsey's theorem.

The Ramsey property



The Canonical Functions Proposition

Proposition (Bodirsky + MP + Tsankov '11)

Let \mathbb{A} be homogeneous ordered Ramsey, and \mathbb{B} be ω -categorical.

Let $f: \mathbb{A} \rightarrow \mathbb{B}$.

Then $\overline{O(f)}$ contains a canonical function $f': \mathbb{A} \rightarrow \mathbb{B}$.

New proof.

For $g, g' \in \overline{O(f)}$ set $g \sim g'$ if $g \in \overline{\{\beta g' \mid \beta \in \text{Aut}(\mathbb{B})\}}$.

Factoring by \sim yields compact space (ω -categoricity).

Consider action of $\text{Aut}(\mathbb{A})$ on that space: $([g]_{\sim}, \alpha) \mapsto [g\alpha^{-1}]_{\sim}$.

By Kechris-Pestov-Todorčević '05, $\text{Aut}(\mathbb{A})$ is **extremely amenable**:

Cont. actions of $\text{Aut}(\mathbb{A})$ on compact Hausdorff spaces have fixed point.

Let $[f']_{\sim}$ be a fixed point of our action.

f' is canonical.

Enjoying canonical functions

In proof of Thomas' theorem:

$$f \in \mathcal{G} \setminus \text{Aut}(V, E)$$

$f: (V, E, a, b) \rightarrow (V, E)$, where $(a, b) \in E$ is sent to a non-edge.

$\overline{O(f)}$ (with respect to action of $\text{Aut}(V, E) \times \text{Aut}(V, E, a, b)$) contains canonical function f' .

- f' is still "contained" in \mathcal{G}
- f' still sends (a, b) to non-edge.

Note:

(V, E) is not Ramsey \implies
work with **Ramsey expansion** $(V, E, <)$.

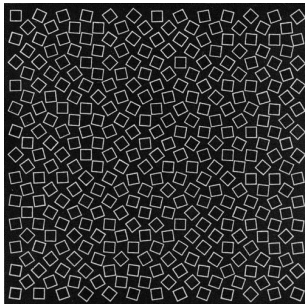
More applications

Reducts:

- Rational numbers $(\mathbb{Q}, <)$ (3) (Cameron '76)
- Random partial order (P, \leq) (3)
(Pach + MP + Pluhár + Pongrácz + Szabó '14)
- Random ordered graph $(V, E, <)$ (42)
(Bodirsky + MP + Pongrácz '15)
- Random digraph (9) (Agarwal '16)
- Homogeneous binary branching C-relation (1)
(Bodirsky + Jonsson + Van Pham '16)
- Homogeneous binary branching semilinear order (1)
(Bodirsky + Bradley-Williams + MP + Pongrácz '17)

Constraint Satisfaction Problems:

- Homogeneous binary branching C-relation
(Bodirsky + Jonsson + Van Pham '16)
- All homogeneous graphs (Bodirsky + Martin + MP + Pongrácz '16)



IV: Canonical functions in the real world

Necessity of the Ramsey property

What if $f: \mathbb{A} \rightarrow \mathbb{B}$, where \mathbb{A} is not Ramsey?

Ramsey property: all continuous actions of $\text{Aut}(\mathbb{A})$ on compact Hausdorff spaces have fixed point.

Canonical functions: fixed points of a particular action.

Problem

Suppose \mathbb{A} is ordered homogeneous, and $\overline{O(f)}$ contains a canonical function for all $f: \mathbb{A} \rightarrow \mathbb{B}$, for all ω -categorical \mathbb{B} .

Is \mathbb{A} Ramsey?

Theorem (Van Pham '17)

Yes, if \mathbb{A} has an ω -categorical Ramsey expansion \mathbb{A}' .

Proof idea. Consider the identity function $id: \mathbb{A} \rightarrow \mathbb{A}'$.

Ramsey expansions

Problem (Hubička + Nešetřil, Bodirsky + MP '11)

Does every ω -categorical structure have an ω -categorical Ramsey expansion?

Theorem (Evans '15)

No.

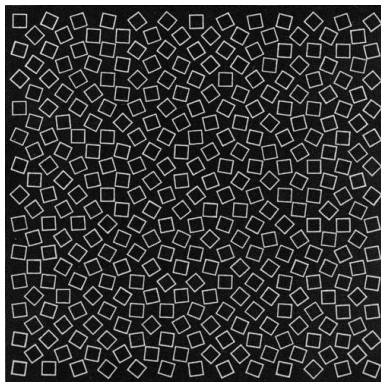
Open for homogeneous structures in finite relational language.

Conjecture (Thomas '91)

Let \mathbb{A} be homogeneous in a finite language.

Then \mathbb{A} has only finitely many reducts.

- Ramsey expansion \implies canonical functions.
- No Ramsey expansion \implies perhaps no canonical functions...



Gerhard von Graevenitz *Regularity - Irregularity V*